Verification in Staged Tile Self-Assembly



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DNA tile self-assembly

DNA tile self-assembly



DNA tile self-assembly



DNA



GAAGTTTGGCGTTAGAACGTTGAAATCCGCCTTGTTAAGACCCCGTCTAAGCA

Single strand













DNA tile self-assembly



DNA tile binary counters [Evans, 2014]





DNA tile



Tile

[Winfree 1996]

(Also Adleman? Rothemund?)



(Also Adleman? Rothemund?)



Tile types





Producible assemblies





Terminal assemblies











Producible assemblies







Terminal assemblies





Tile types





Producible assemblies





Terminal assemblies





Tile types





Producible assemblies







Terminal assemblies



Two-handed tile assembly model (2HAM)



This paper is about **verifying** the **behavior** of 2HAM systems.

Verification problems



maybe some other stuff

<u>Output</u>

"Yes"

(system has a specific behavior)

or

"No"

(system does not have that behavior)

Verification problems



Specific verification problems

2HAM Unique Assembly Verification (UAV) Problem



<u>Output</u>

"Yes"

(+ + + is the unique terminal assembly of the input system)

or

"No"

 is not a terminal assembly or there is another terminal assembly of the input system)

2HAM Unique Assembly Verification (UAV) Problem




<u>Output</u>

"Yes"

(+ + + is the unique terminal assembly of the input system)

or

"No"

 is not a terminal assembly or there is another terminal assembly of the input system)



<u>Output</u>



(+ + + is the unique terminal assembly of the input system)

or







"Yes"

(is the unique terminal assembly of the input system)

or

"No"

(is not a terminal assembly or there is another terminal assembly of the input system)



Assembly





Assembly

<u>Output</u>



(is the unique terminal assembly of the input system)

or

"No"

there is another terminal assembly of the input system)



<u>Output</u>



assembly of the input system)

or

"No"

(there is a terminal assembly of the input system without shape





"No"

(there is a terminal assembly of the input system without shape



Shape

<u>Output</u>



assembly of the input system)

or

"No"

(there is a terminal assembly of the input system without shape





<u>Output</u>



is the shape of every terminal assembly of the input system)

or

"No"

(there is a terminal assembly of the input system without shape)



Our first result

Complexity classes





<u>NP</u>

"Yes" can be verified in polynomial time.

Returns "Yes" if any branch returns "Yes".

NP-complete: does a Boolean formula have satisfying assignment?

<u>coNP</u>

"No" can be verified in polynomial time.

Returns "No" if any branch returns "No".

coNP-complete: Is every assignment of a Boolean formula satisfying?

1. 2HAM USV is coNP^{NP}-complete.

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An oracle

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An oracle

The coNP algorithm can make O(1)-time calls to any NP algorithm.

1. 2HAM USV is coNP^{NP}-complete.



An *oracle* ----

The coNP algorithm can make O(1)-time calls to any NP algorithm.

Already known: 2HAM UAV is coNP-complete. [Unpublished]

2HAM USV is coNP^{NP}-hard Reduction is from ∀∃SAT:

Input:

A boolean formula with variables $x_1, x_2, ..., x_j, y_1, y_2, ..., y_k$

Output:

Does every assignment of x₁, x₂, ..., x_j have an assignment of y₁, y₂, ..., y_k that satisfies the formula?

∀∃SAT proved coNP^{NP}-hard by [Stockmeyer 1975]

2HAM USV is coNP^{NP}-hard

(The assemblies involved in the reduction)



Assignment assembly

2HAM USV is coNPNP-hard

(The assemblies involved in the reduction)

Satisfying assignment







2HAM USV is coNPNP-hard

(The assemblies involved in the reduction)





2HAM USV is coNPNP-hard

(The assemblies involved in the reduction)



Does every assignment of x₁, x₂ have an assignment of y_1 , y_2 , that satisfies the formula?





x1, x2 assignments















































Filler assembly







































































Reduction from to ∀∃SAT to 2HAM USV

Terminal assemblies



Reduction from to ∀∃SAT to 2HAM USV

Terminal assemblies


Reduction from to ∀∃SAT to 2HAM USV

Terminal assemblies



Reduction from to ∀∃SAT to 2HAM USV

Terminal assemblies





Staged tile self-assembly model









Staged tile assembly model [Demaine et al. 2008]



Staged tile assembly model [Demaine et al. 2008]



Our second and third results

Staged UAV Problem



Staged USV Problem





Shape

<u>Output</u>



or

"No"



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- 2. Staged UAV and USV are coNP^{NP}-hard.

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- 2. Staged UAV and USV are coNP^{NP}-hard.
- 3. Staged UAV and USV are in PSPACE.

Actually, we prove something stronger...

PSPACE













Assembly





- 1. 2HAM USV is coNP^{NP}-complete.
- 2. Staged UAV and USV are coNP^{NP}-hard.
- 3. Staged UAV and USV are in PSPACE. 3.5 For O(1) stages, in $coNP^{NP^{NP} \cdots NP}$

Are staged UAV and USV PSPACE-hard? For O(1) stages, coNPNPNP...NP-hard?

Are staged UAV and USV PSPACE-hard? For O(1) stages, coNPNPNP--NP-hard?

Do results change if only allowed?

Are staged UAV and USV PSPACE-hard? For O(1) stages, coNPNPNP--NP -hard?

Do results change if only allowed?

Are there other self-assembly problems with complexity coNPNPNP...NP or NPNPNP...NP??

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Abstract

We prove the unique assembly and unique shape verification problems, benchmark measures of self-assembly model power, are $coNP^{NP}$ -hard and contained in PSPACE (and in Π_{2s}^{P} for staged systems with s stages). En route, we prove that unique shape verification problem in the 2HAM is $coNP^{NP}$ -complete.

1 Introduction

Here we consider the complexity of two standard problems in tile self-assembly: deciding whether a system uniquely assembles a given assembly or shape. These so-called *unique* assembly and *unique shape verification* problems are benchmark problems in tile assembly, and have been studied in a variety of models, including the aTAM [1, 2], the q-tile model [6], and the 2HAM [3].

The unique assembly and unique shape verification problems ask whether a system behaves as expected: does a given system yield a unique given assembly or assemblies of a given unique shape? The distinct rules by which assemblies form in various tile assembly models