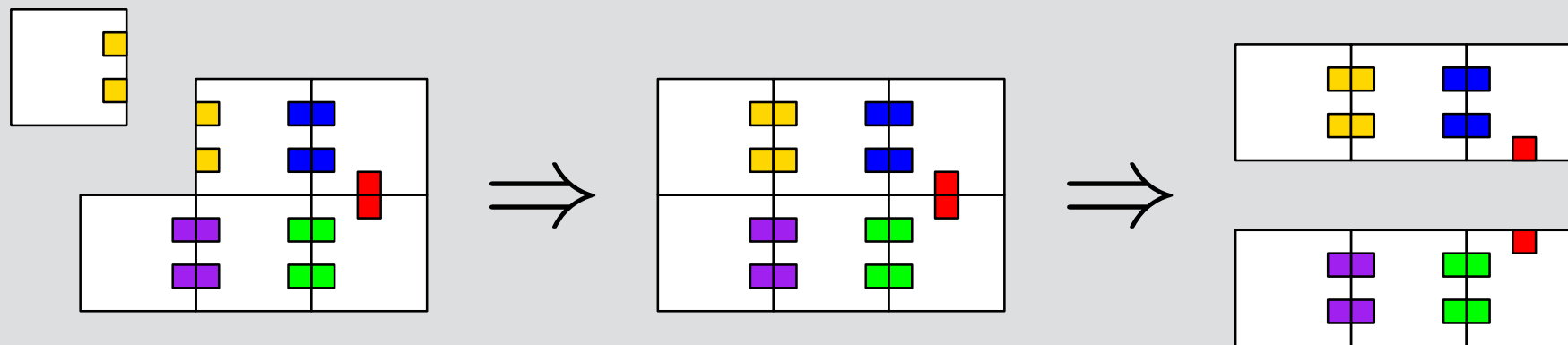


# Size-Dependent Tile Self-Assembly: Constant-Height Rectangles and Stability



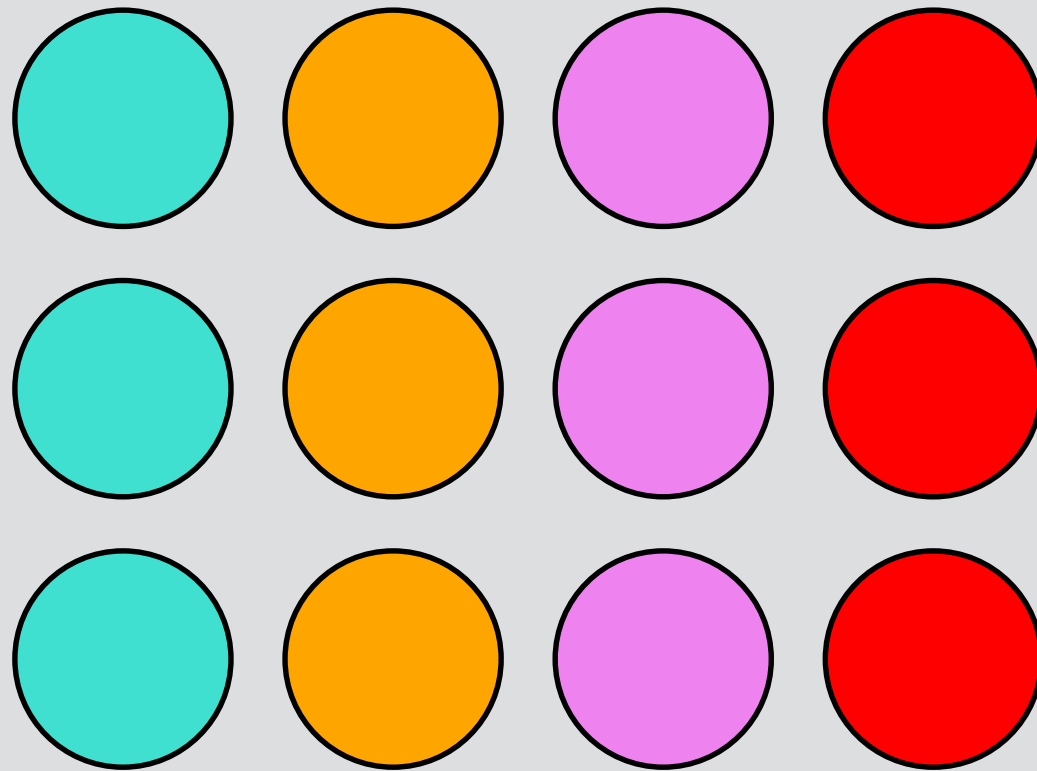
Sándor Fekete, Robert Schweller, Andrew Winslow



# Self-Assembly

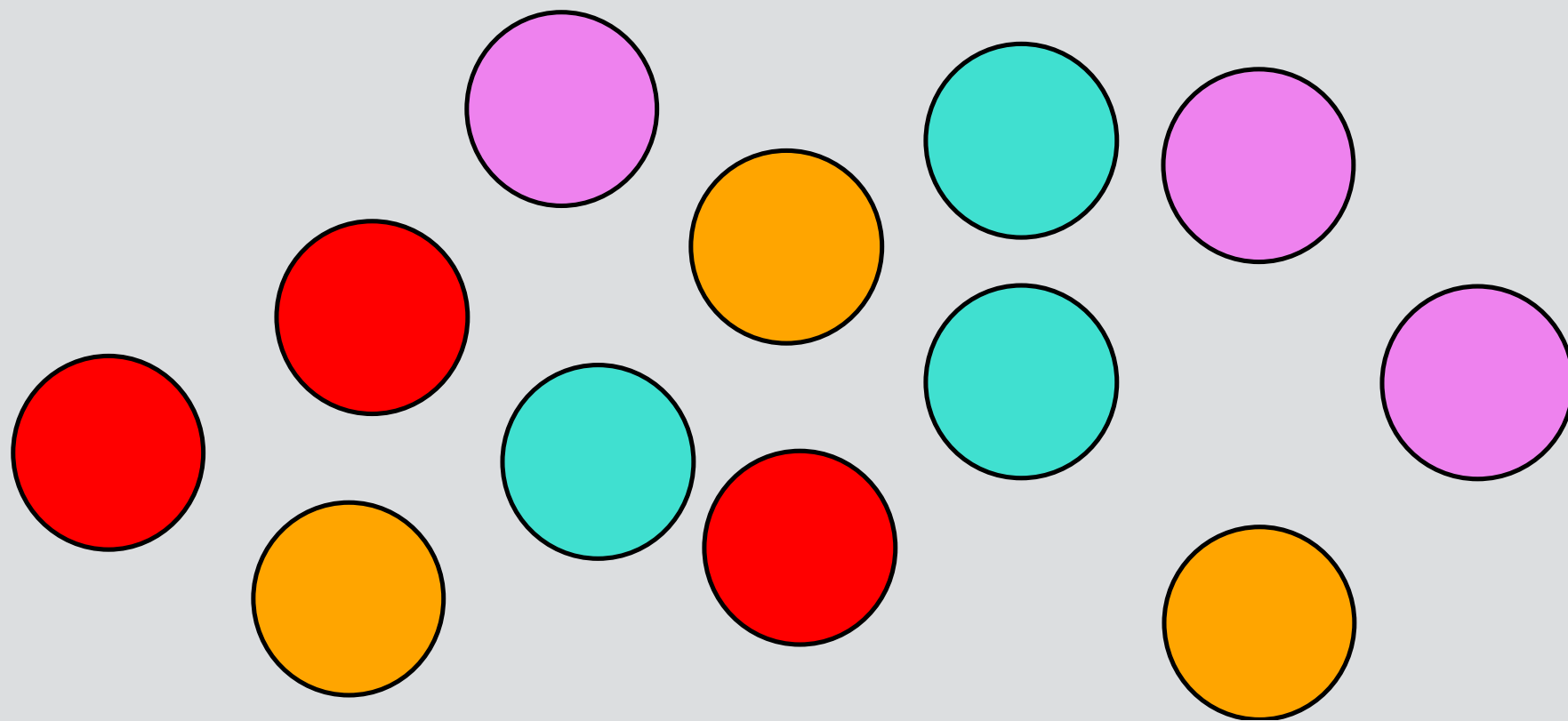
# Self-Assembly

Simple particles coalescing into  
complex superstructures.



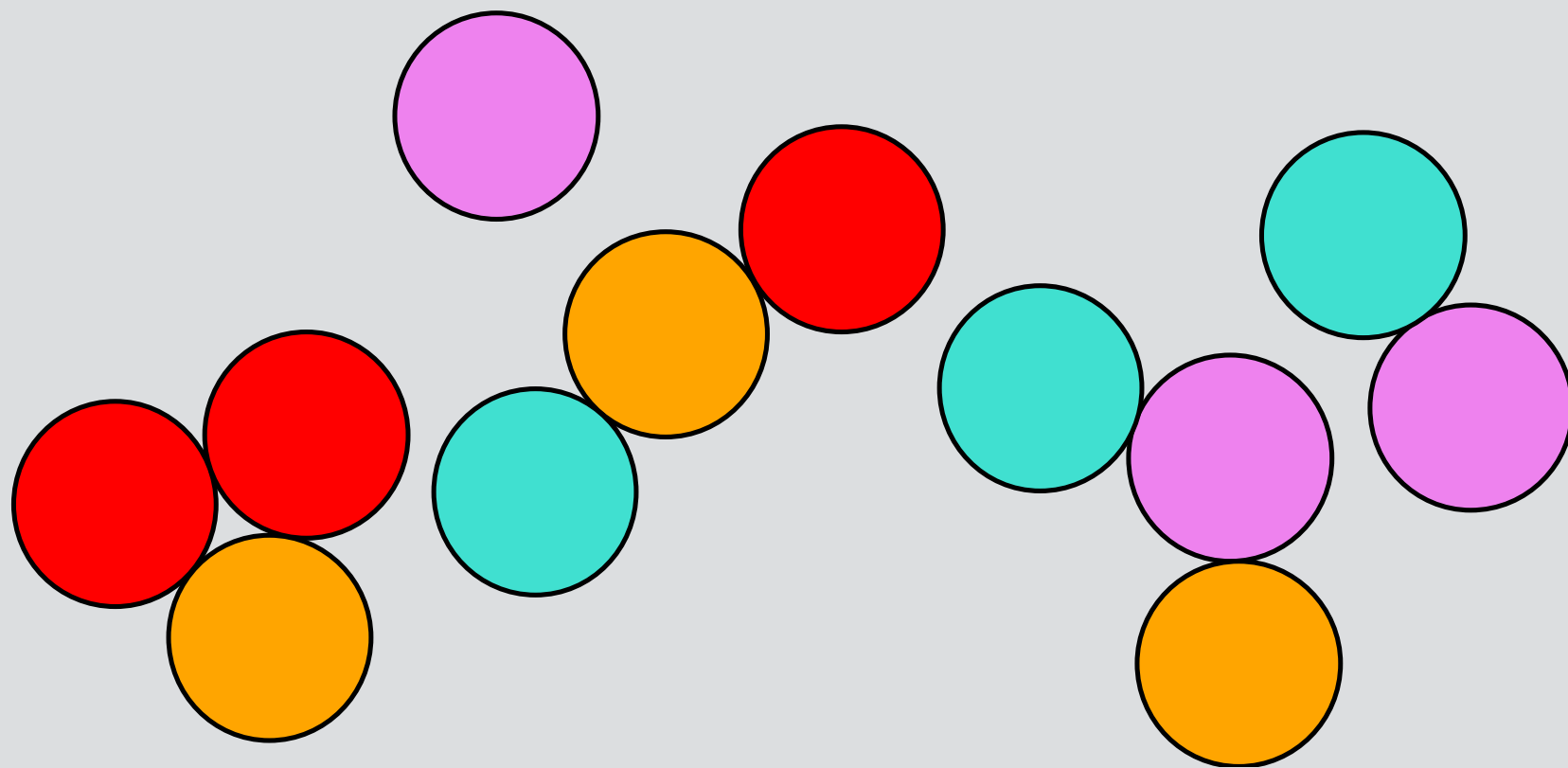
# Self-Assembly

Simple particles coalescing into  
complex superstructures.



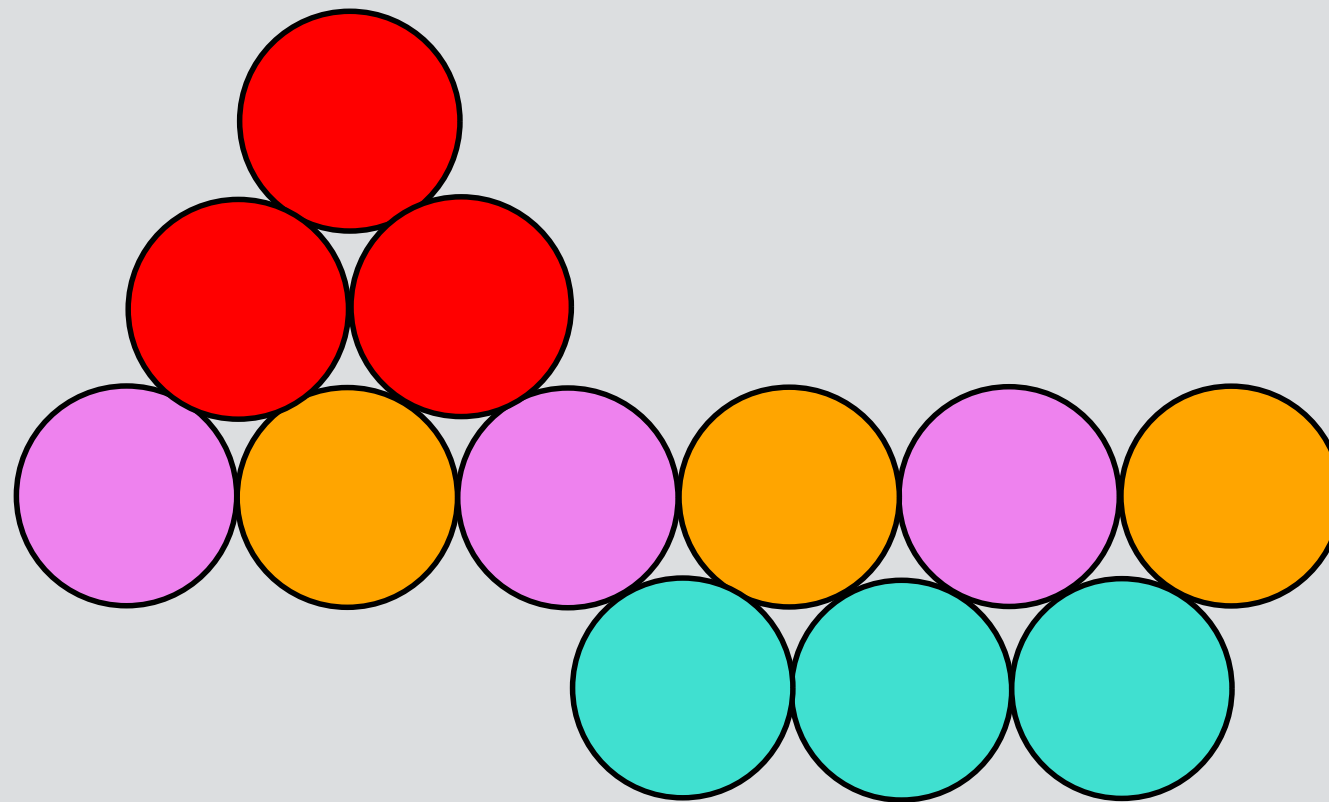
# Self-Assembly

Simple particles coalescing into  
complex superstructures.



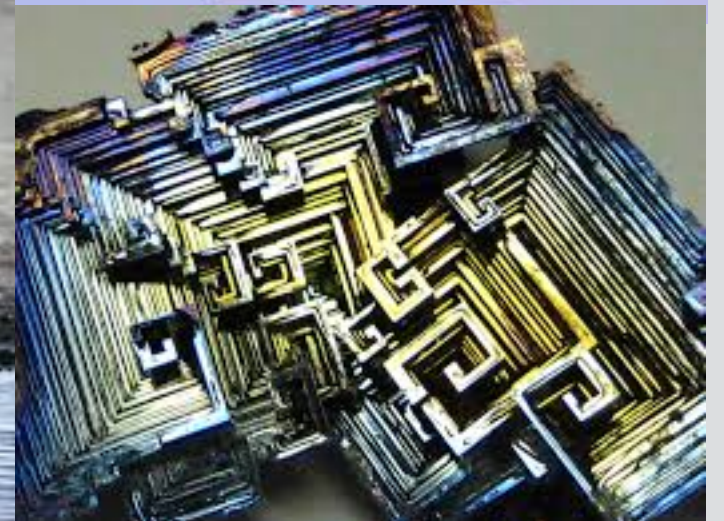
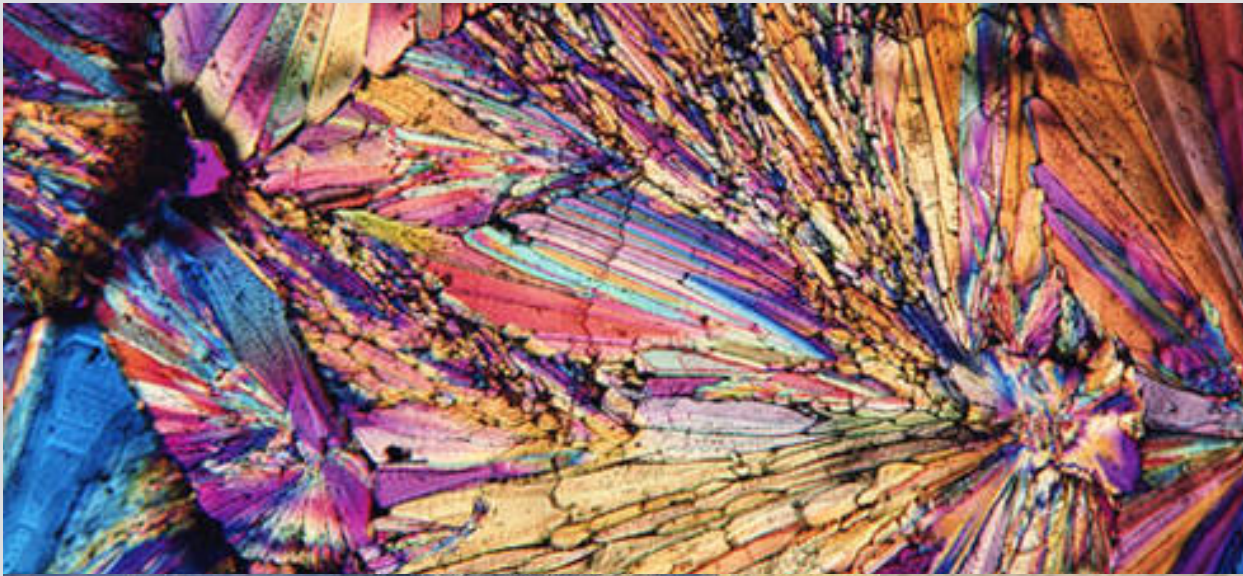
# Self-Assembly

Simple particles coalescing into complex superstructures.



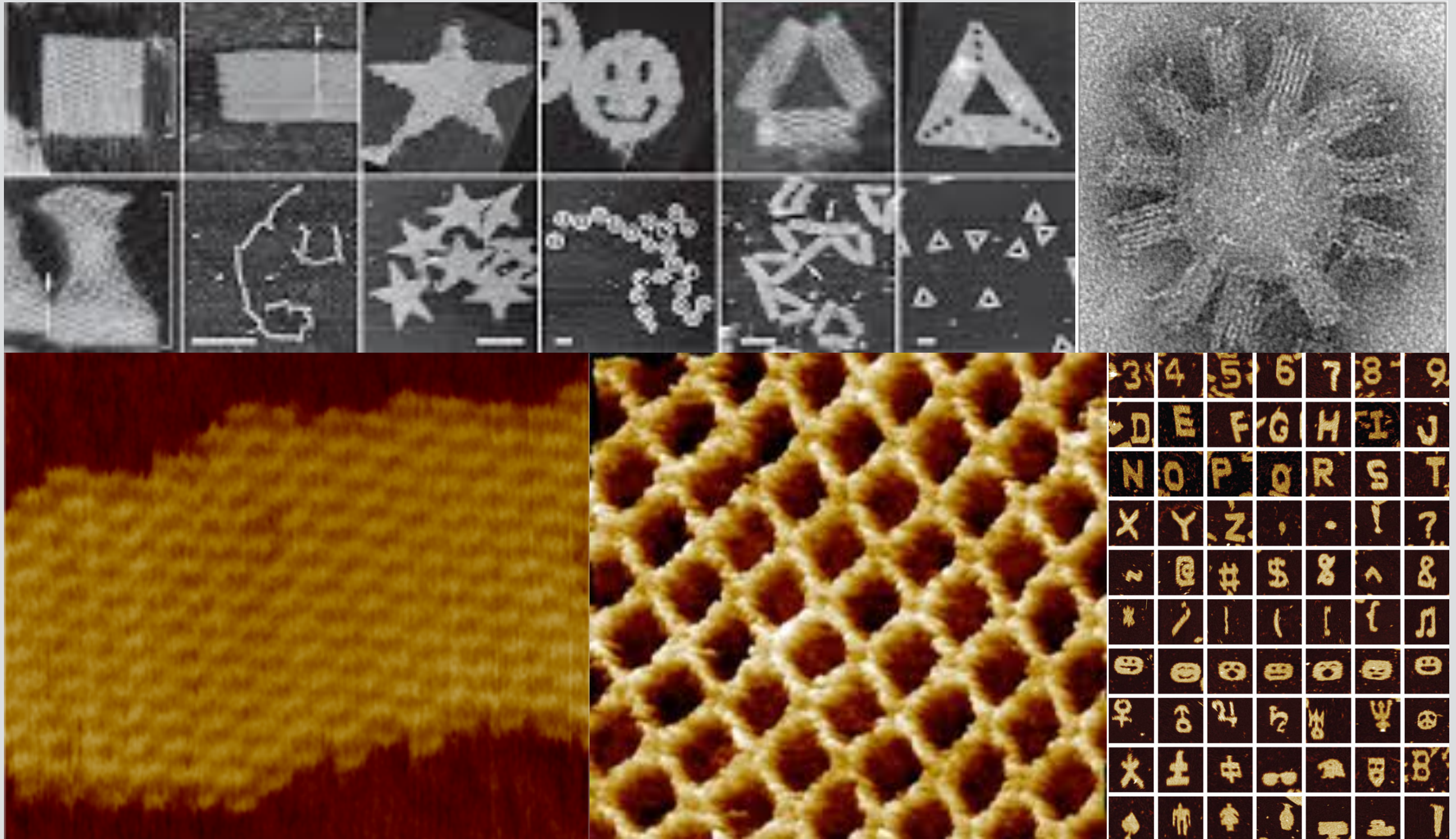


# Natural Self-Assembly



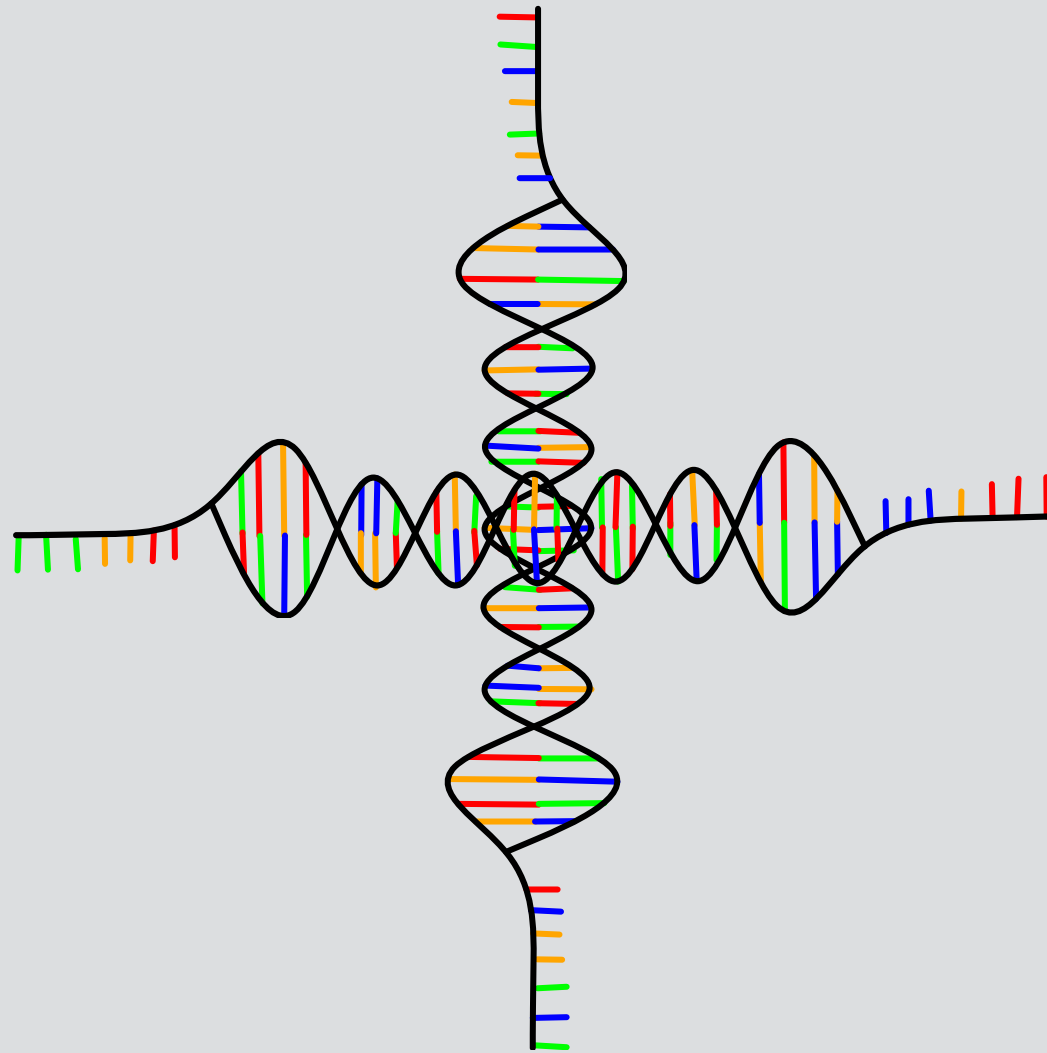


# Synthetic Self-Assembly with DNA

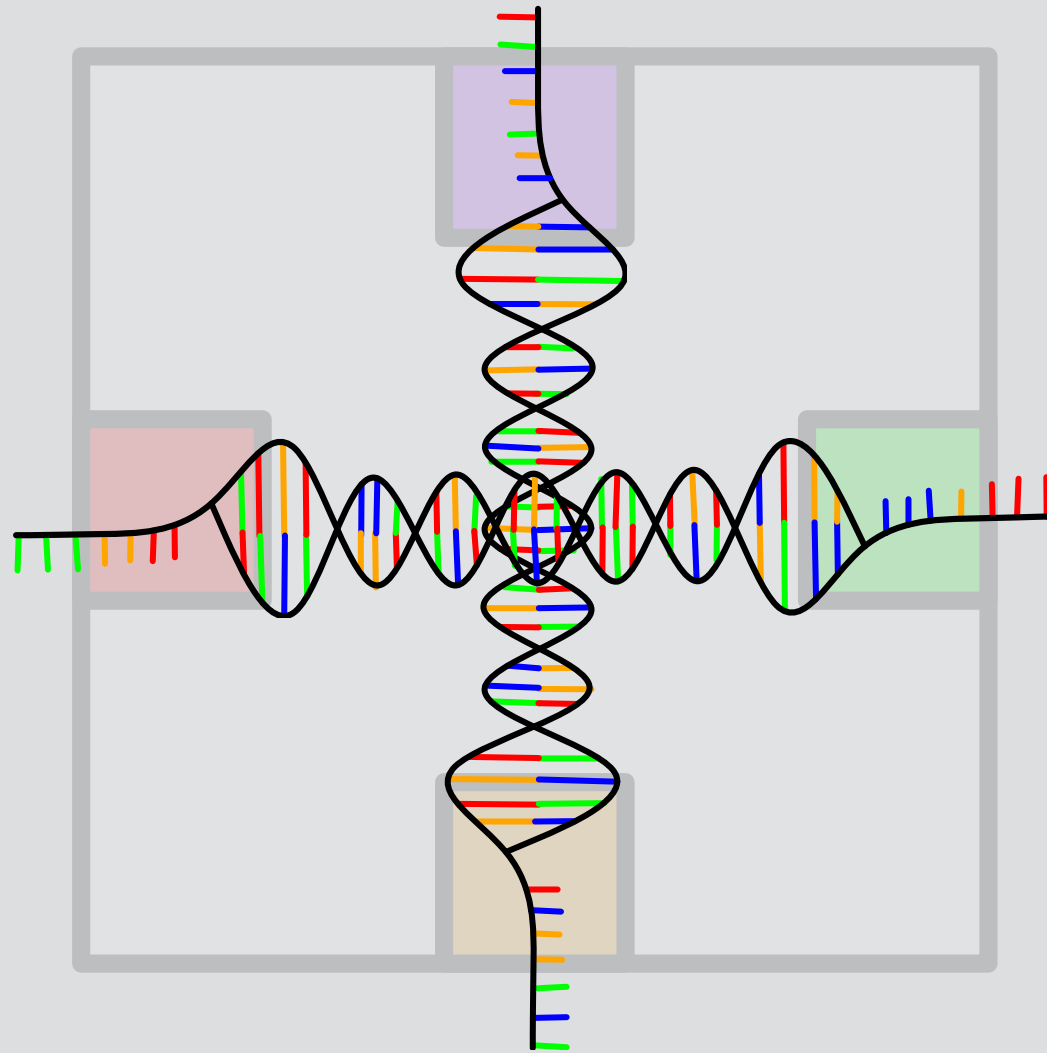




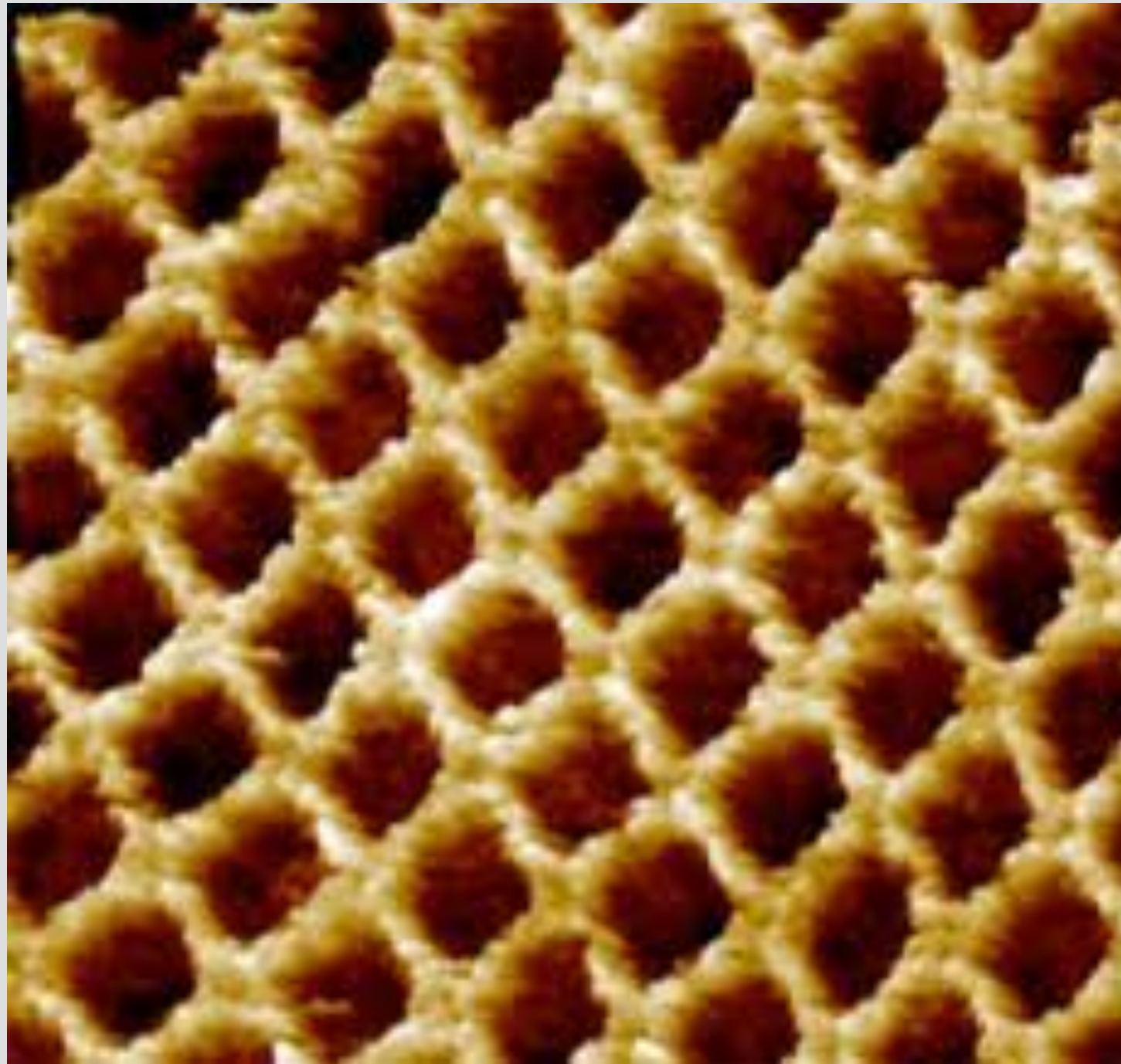
# Synthetic Self-Assembly with DNA



# Synthetic Self-Assembly with DNA

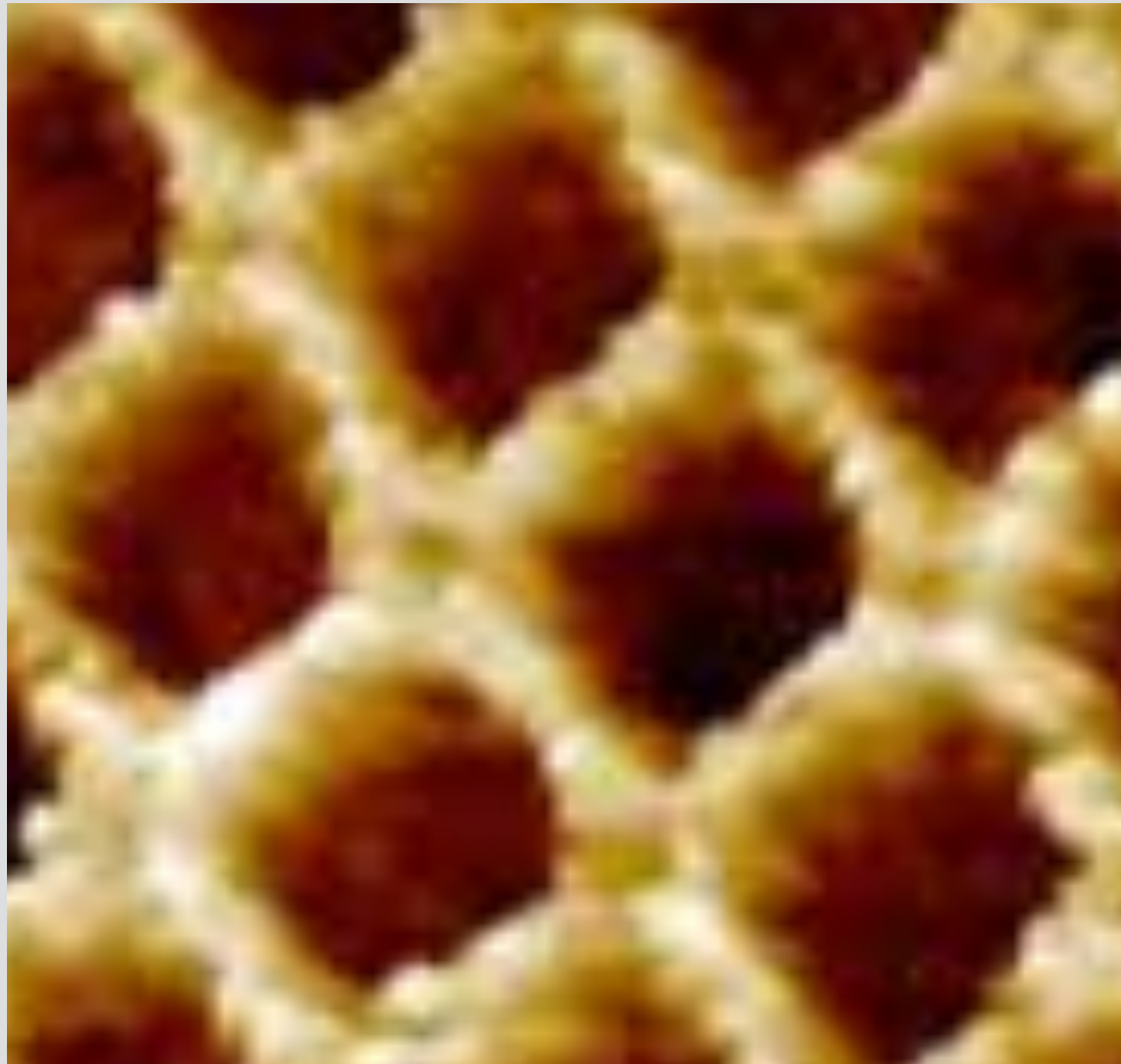


# Synthetic Self-Assembly with DNA

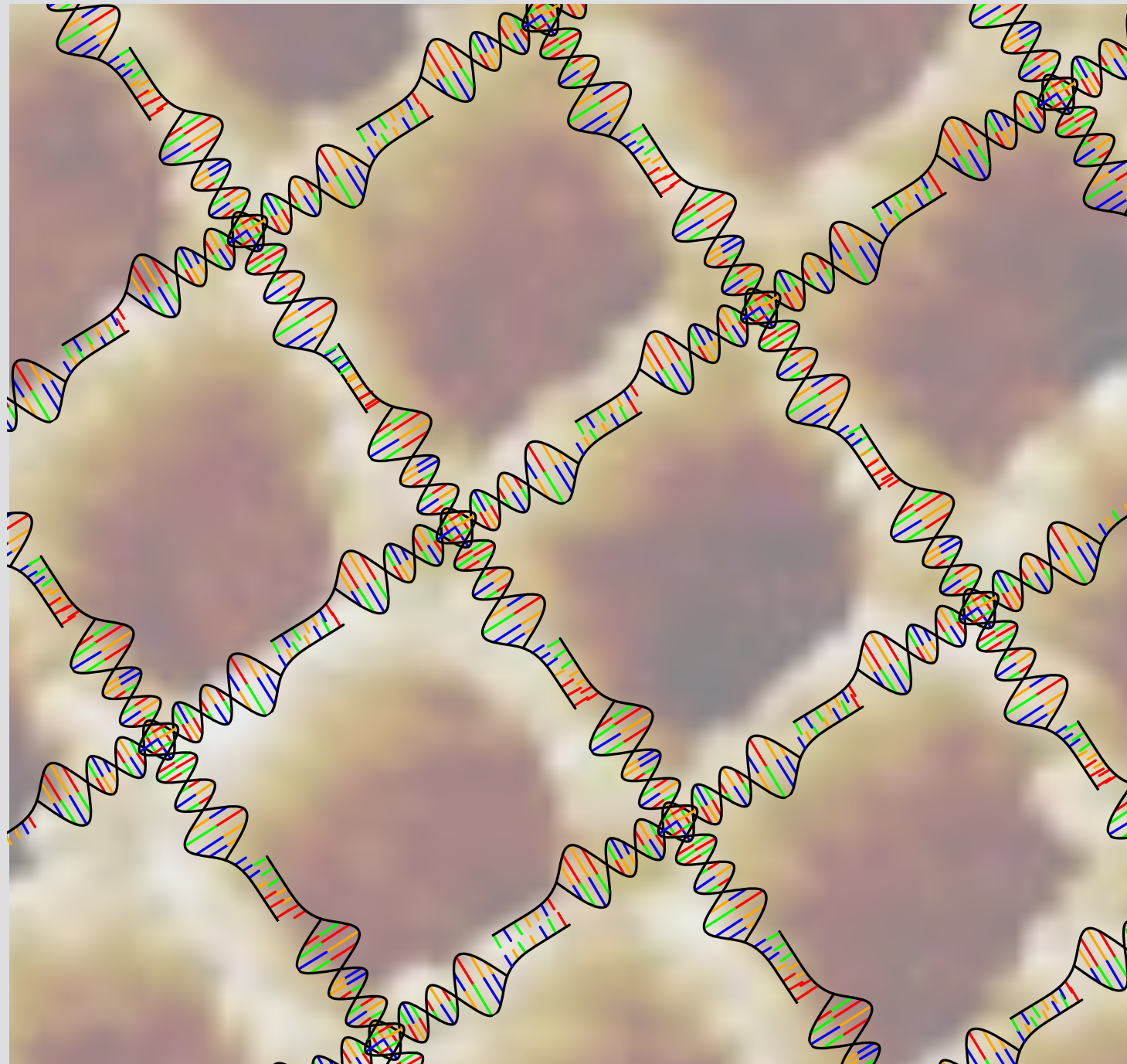




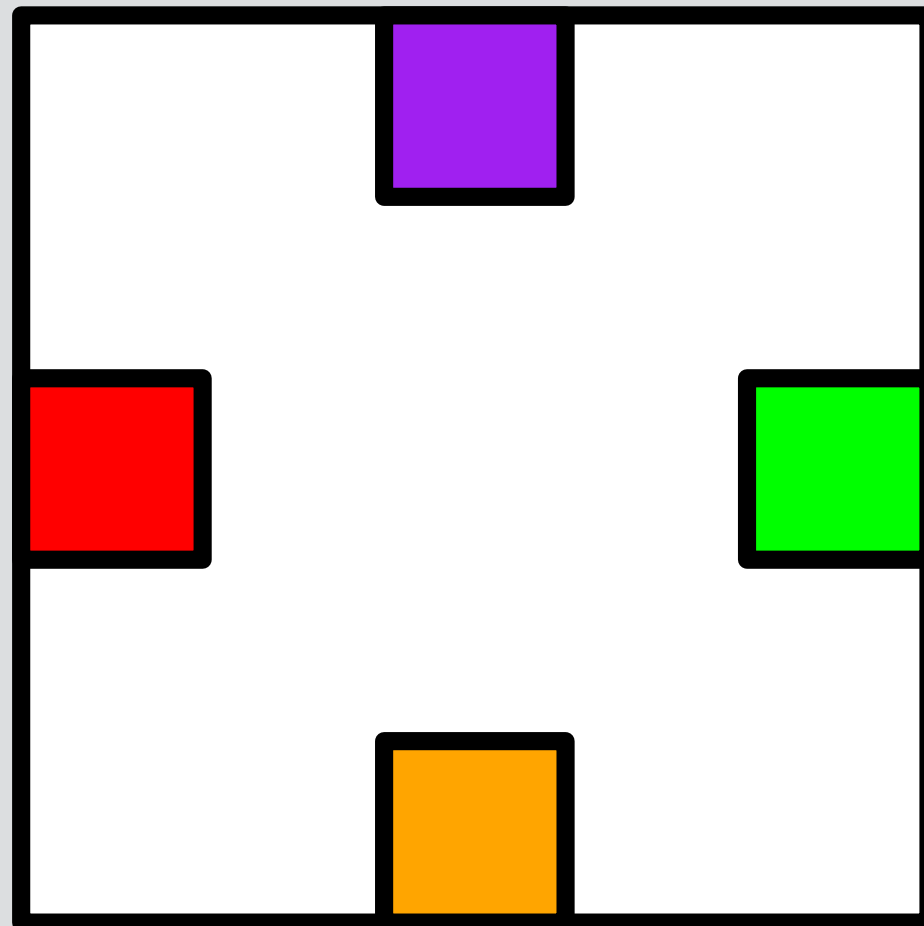
# Synthetic Self-Assembly with DNA



# Synthetic Self-Assembly with DNA

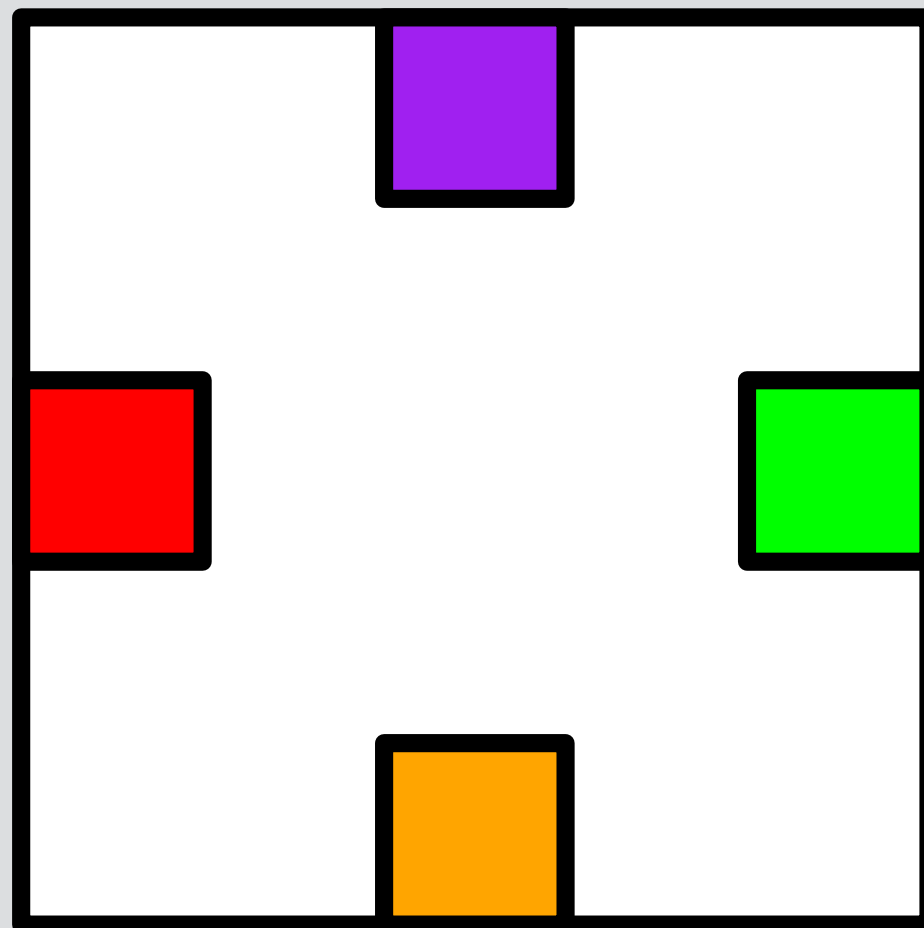


# Tile Self-Assembly





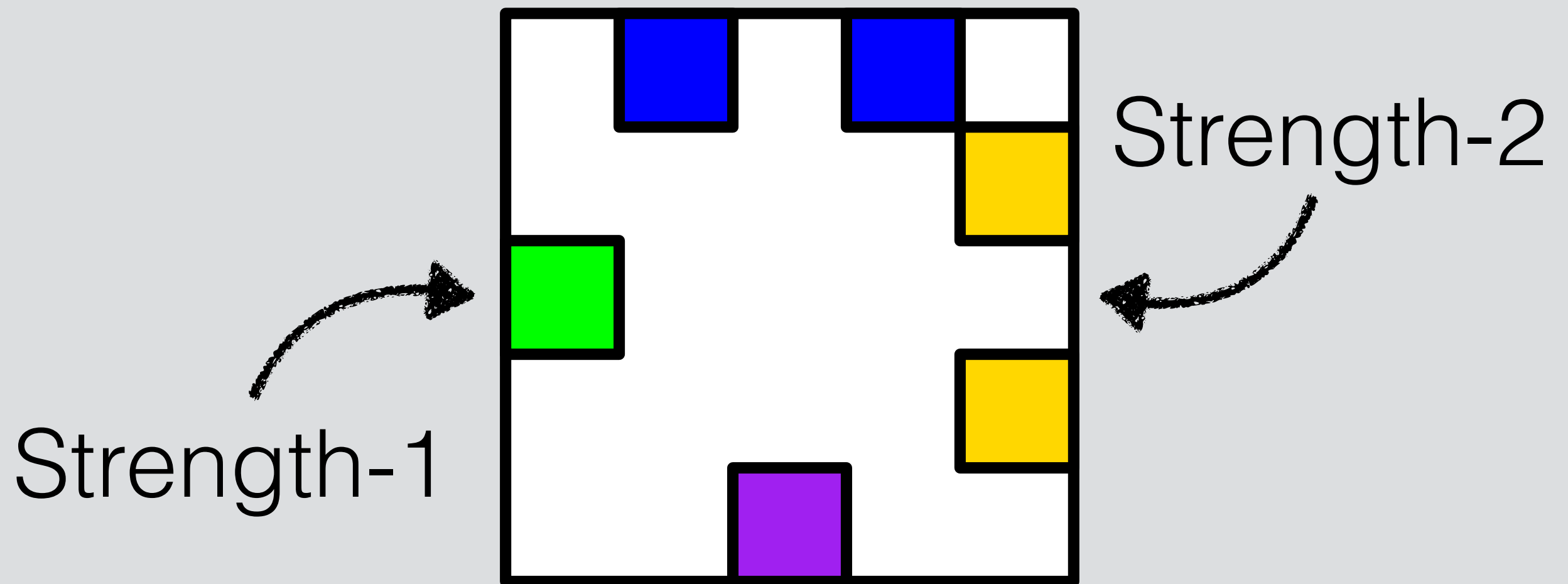
# Tile Self-Assembly



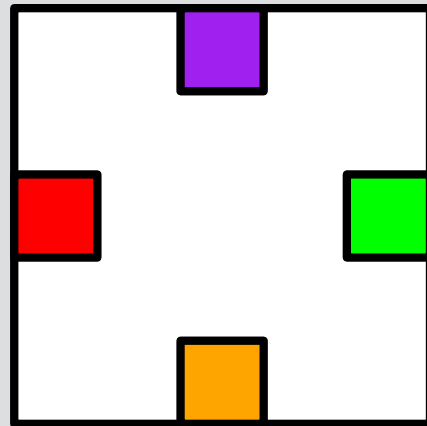
Glues

Two curved black arrows originate from the word 'Glues'. One arrow points to the green tile at (2,3), and the other points to the orange tile at (3,2).

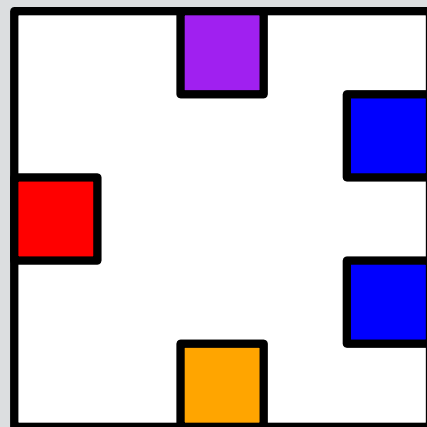
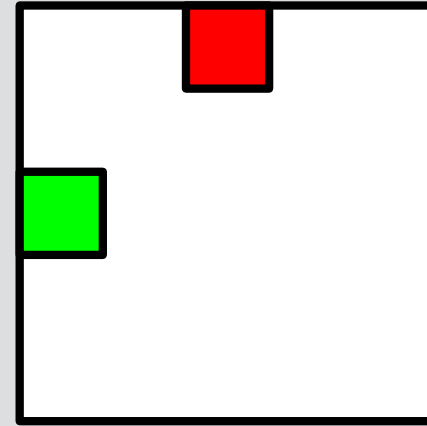
# Tile Self-Assembly



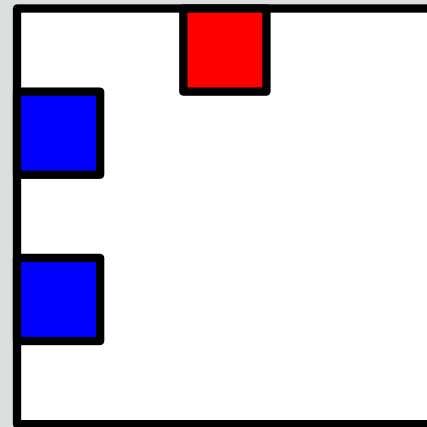
At temperature  $\tau = 1$



$$1 \geq \tau$$

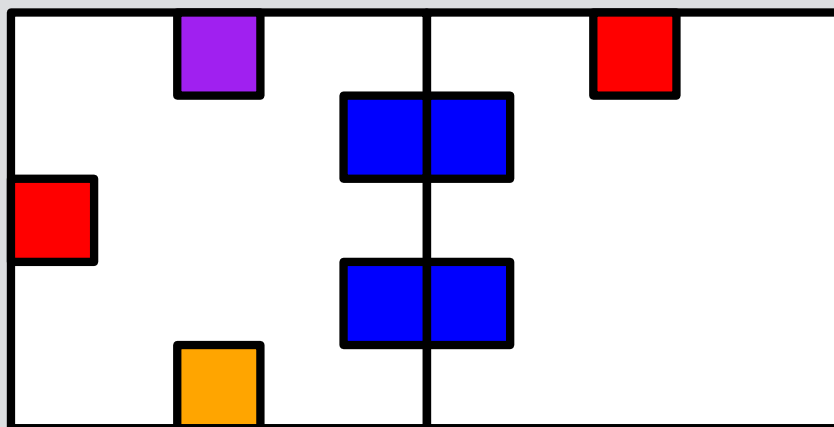
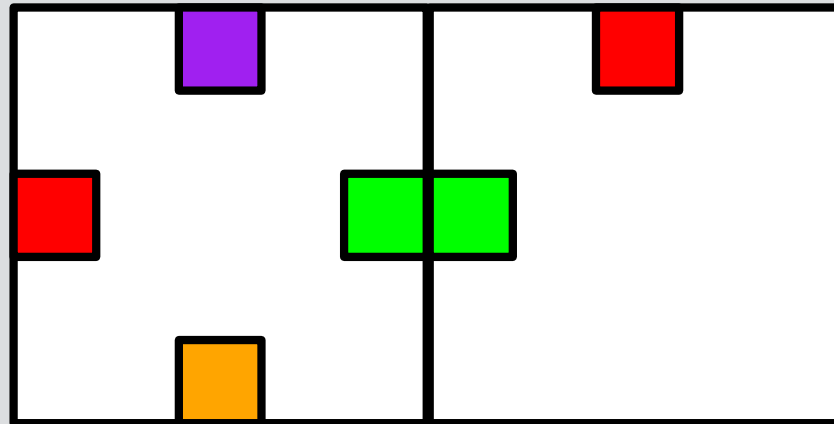


$$2 \geq \tau$$

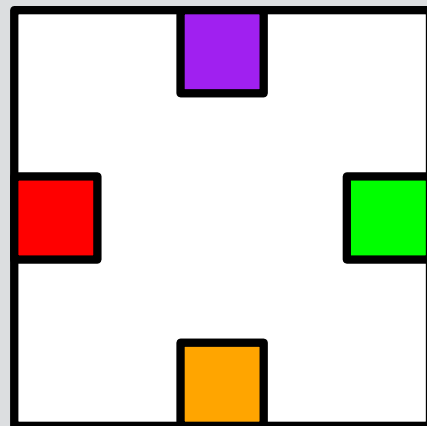




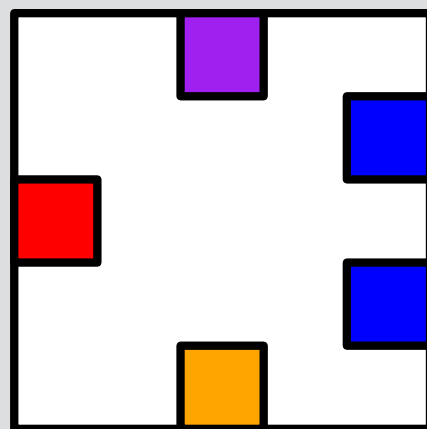
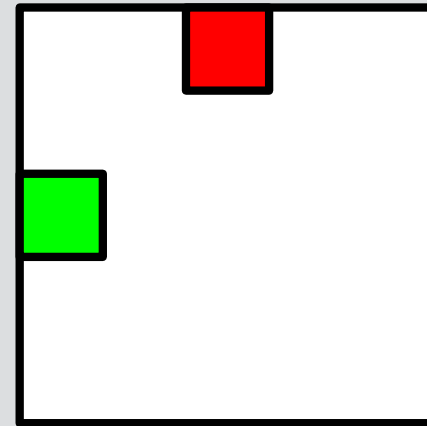
At temperature  $\tau = 1$



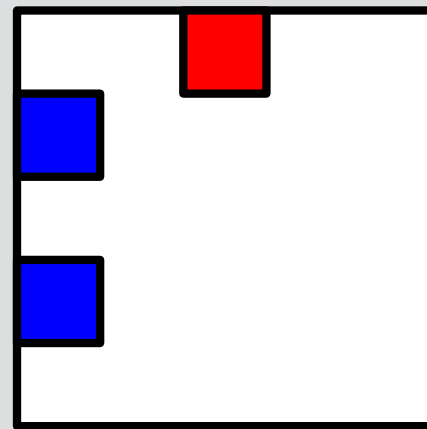
At temperature  $\tau = 2$



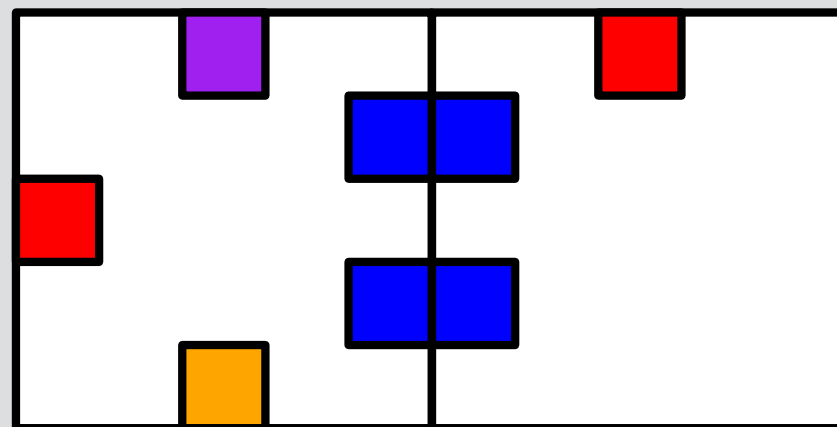
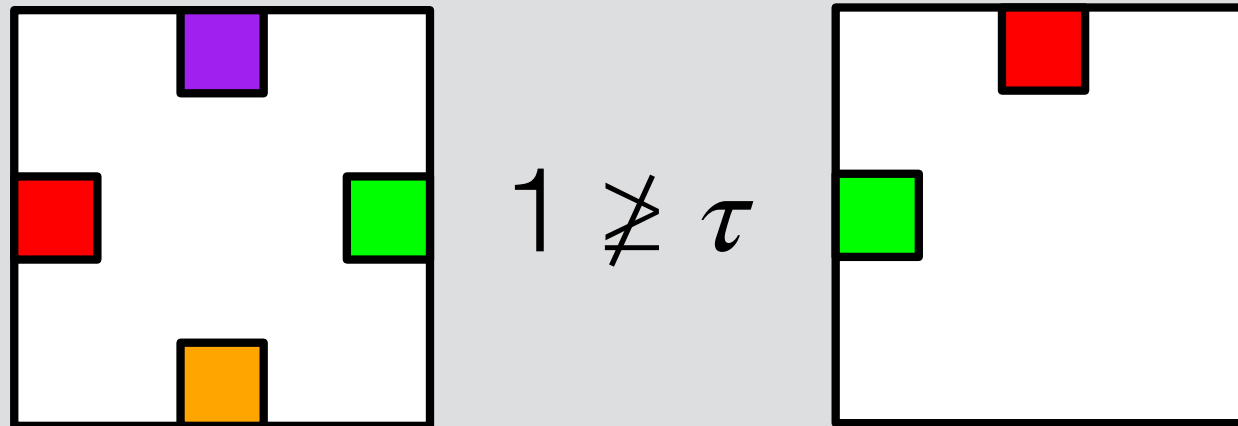
$1 \not\geq \tau$



$2 \geq \tau$

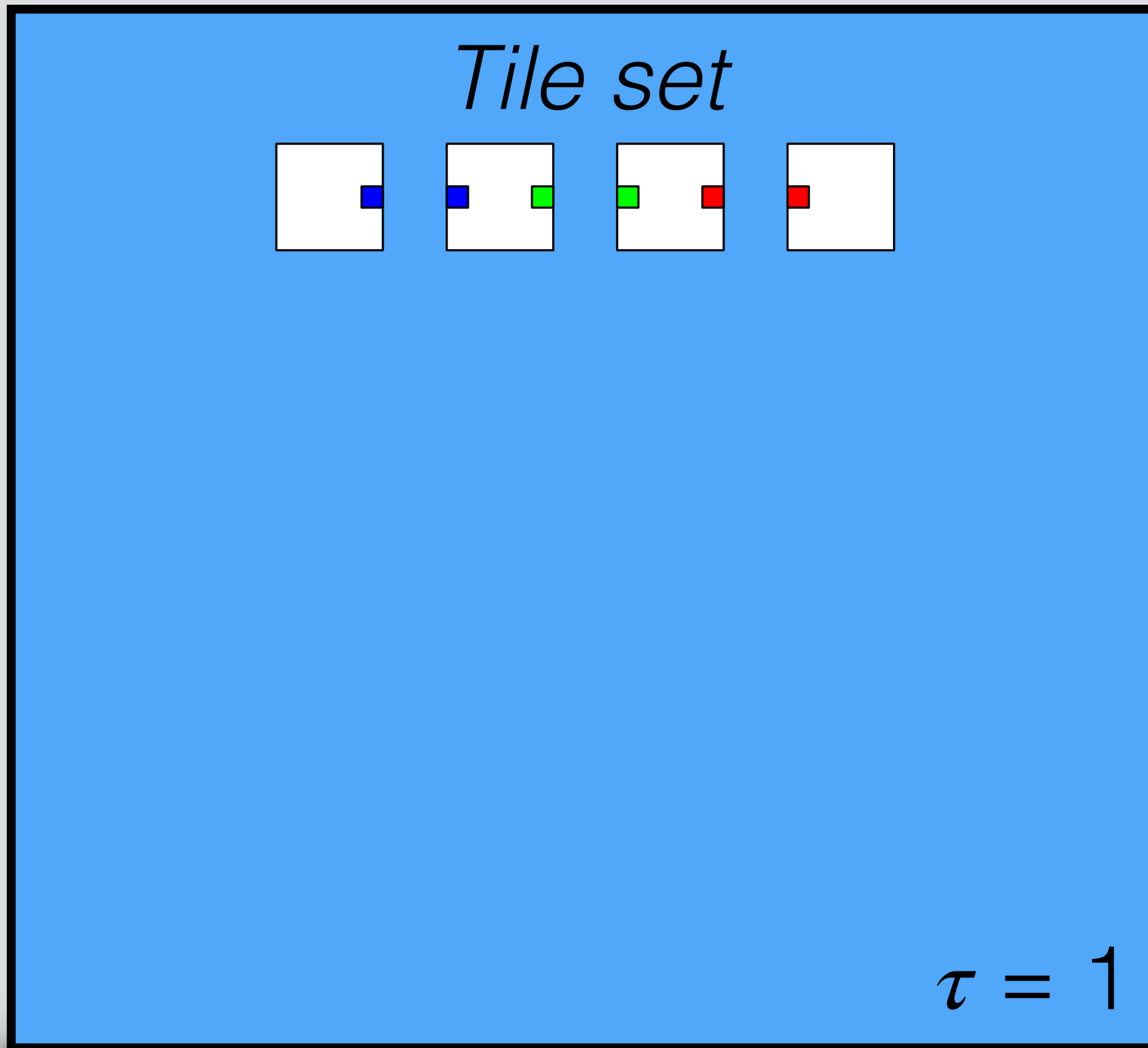


At temperature  $\tau = 2$



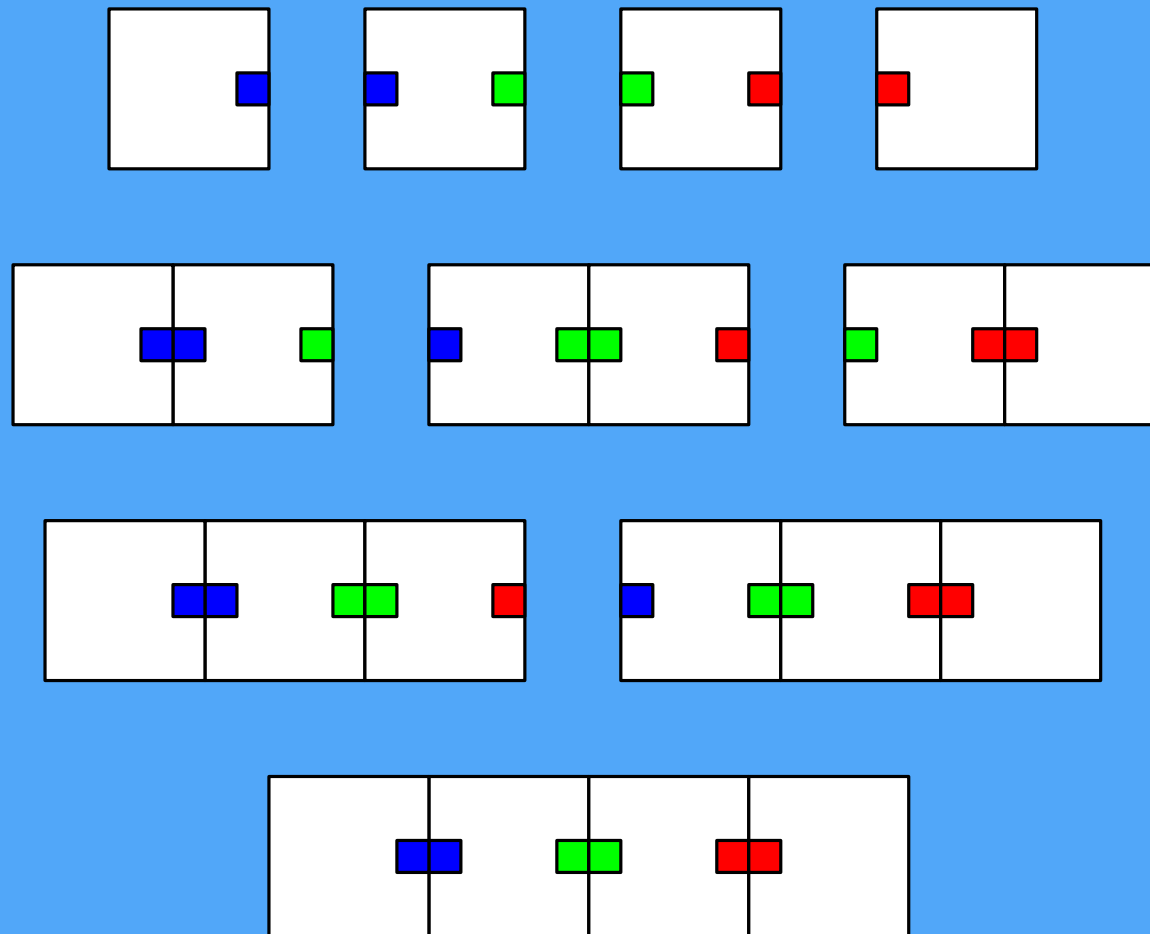


# Two-handed assembly



# Two-handed assembly

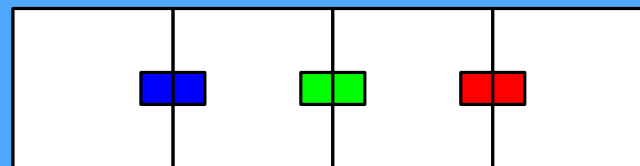
*Producible assemblies*



$$\tau = 1$$

# Two-handed assembly

*Terminal assemblies*



$$\tau = 1$$



$\tau = 1$



$\tau = 2$

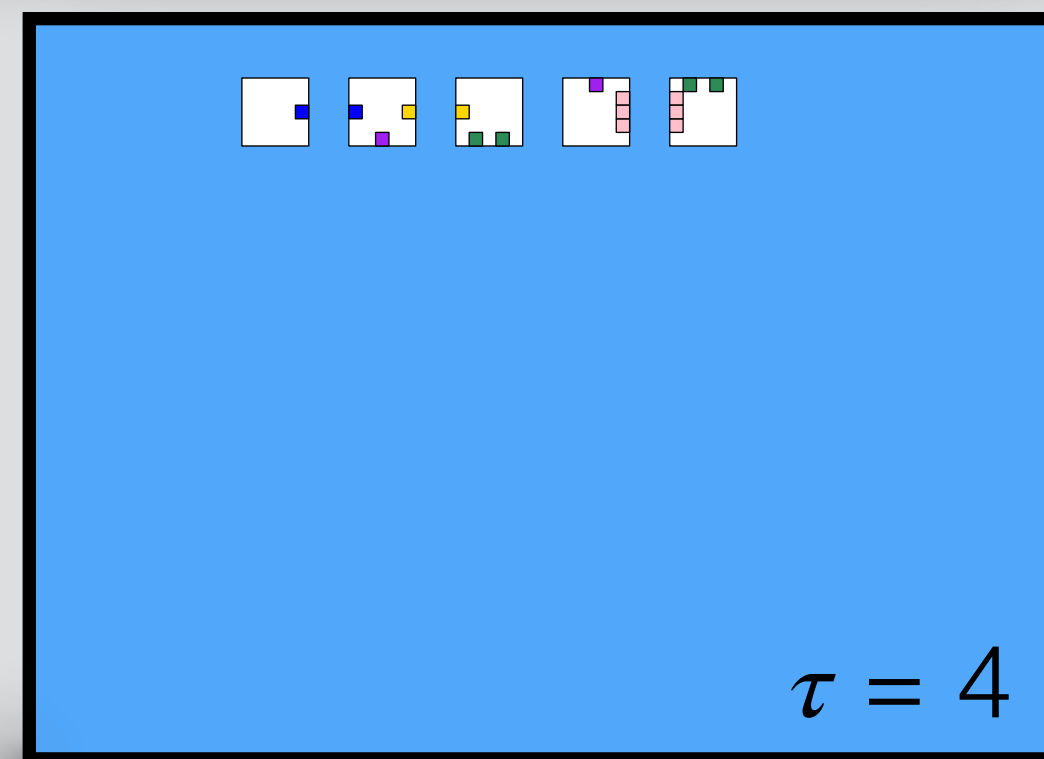
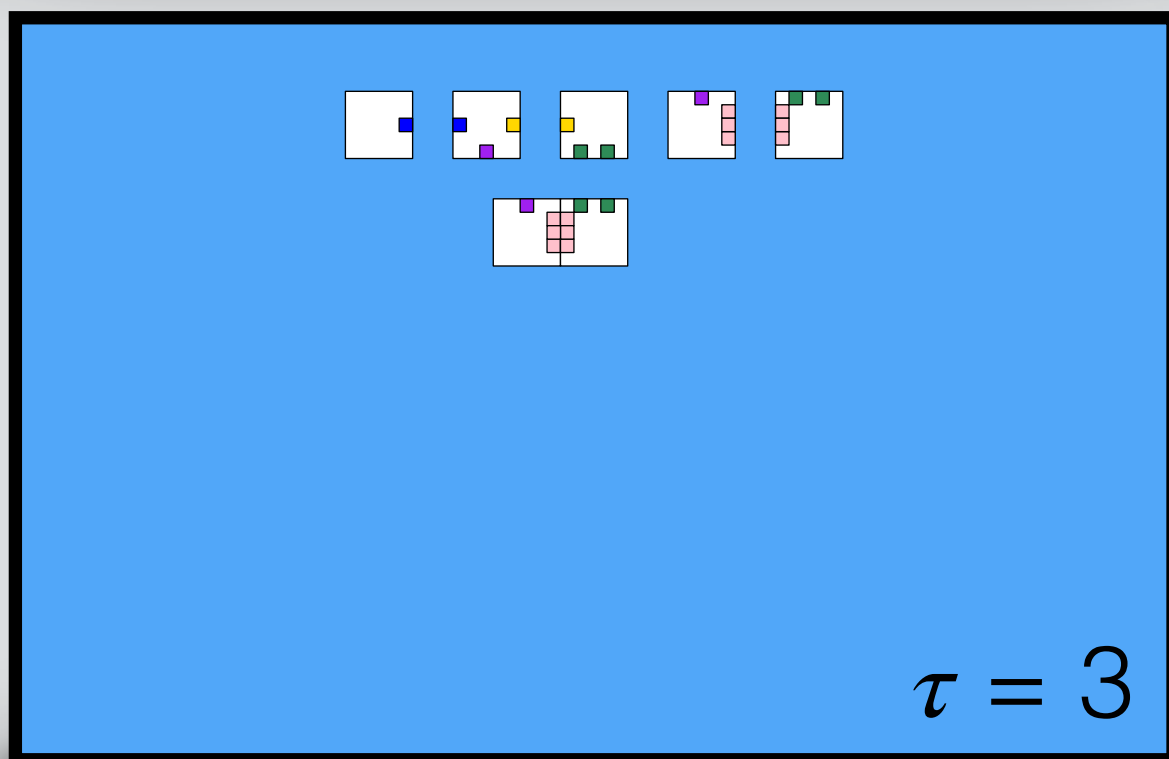
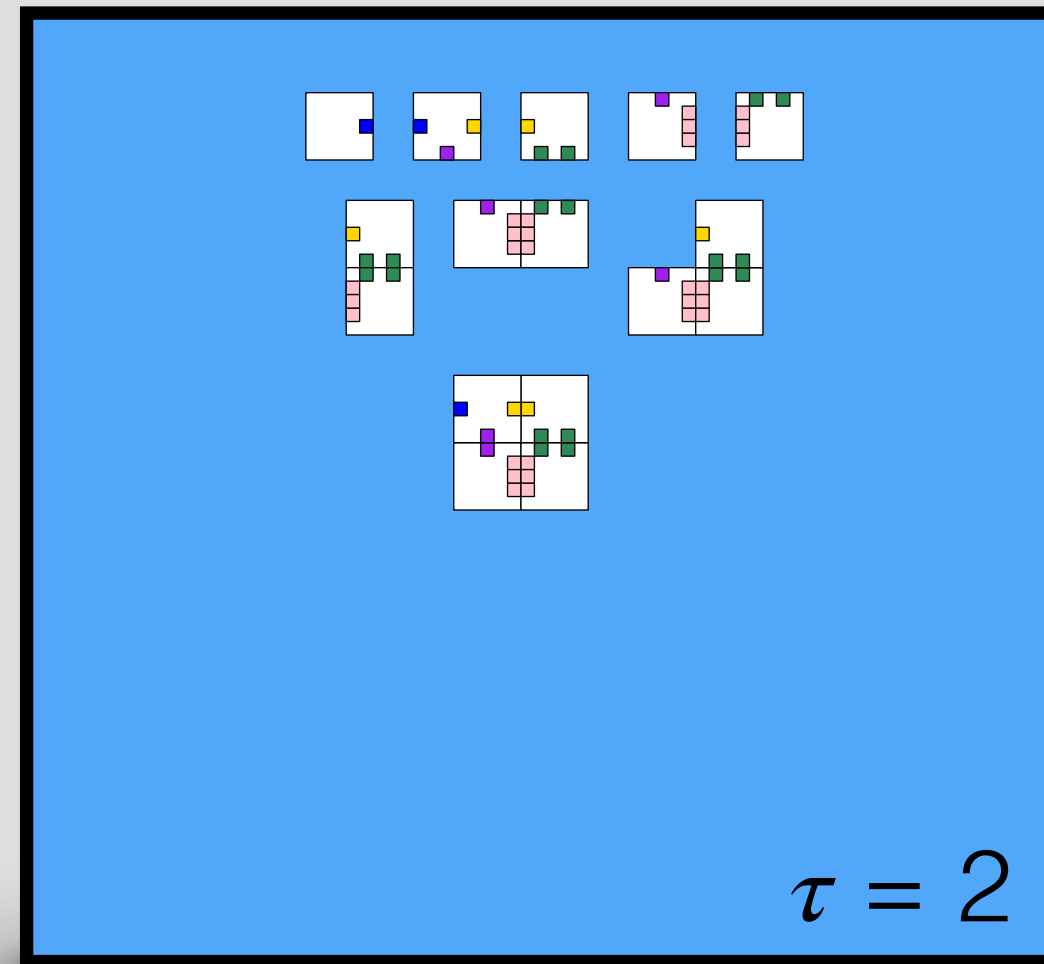
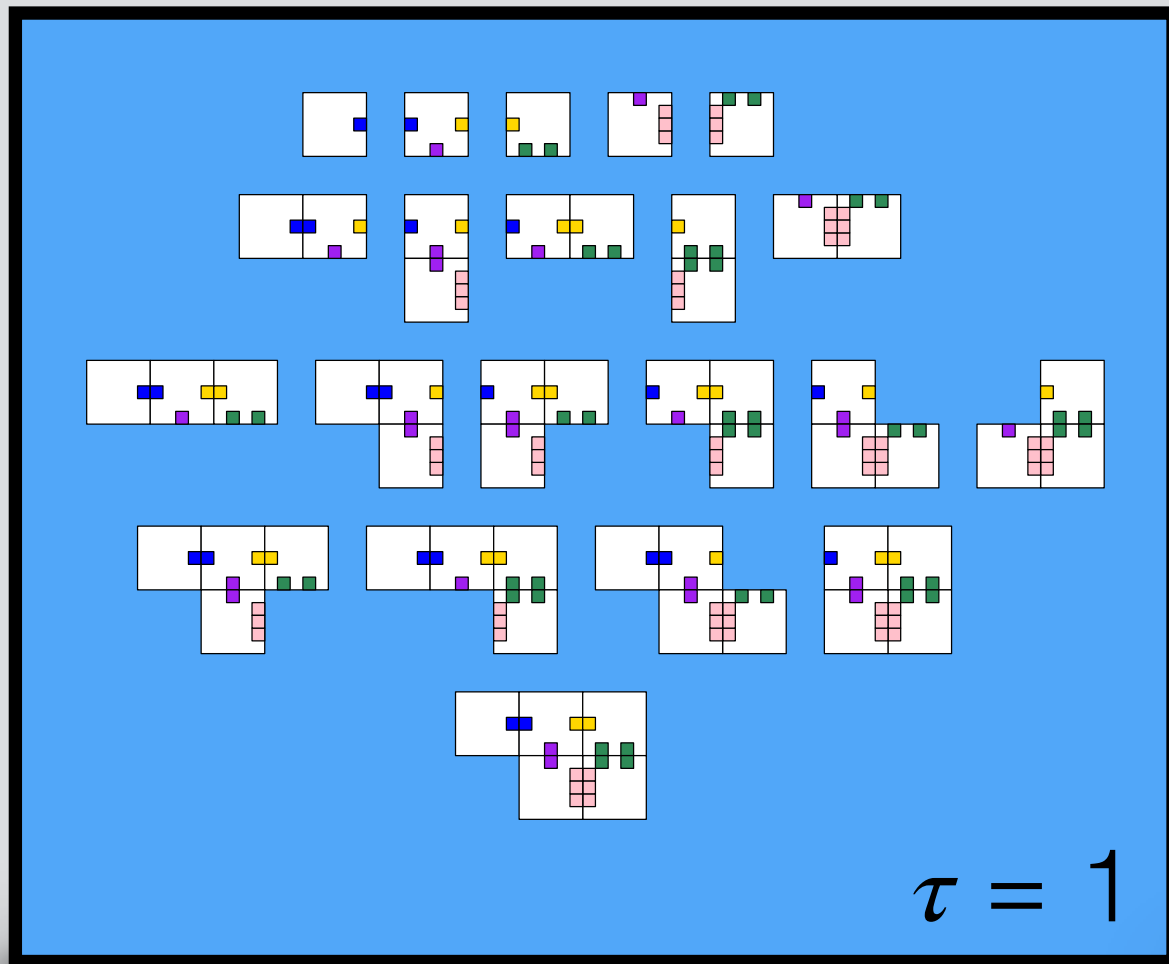


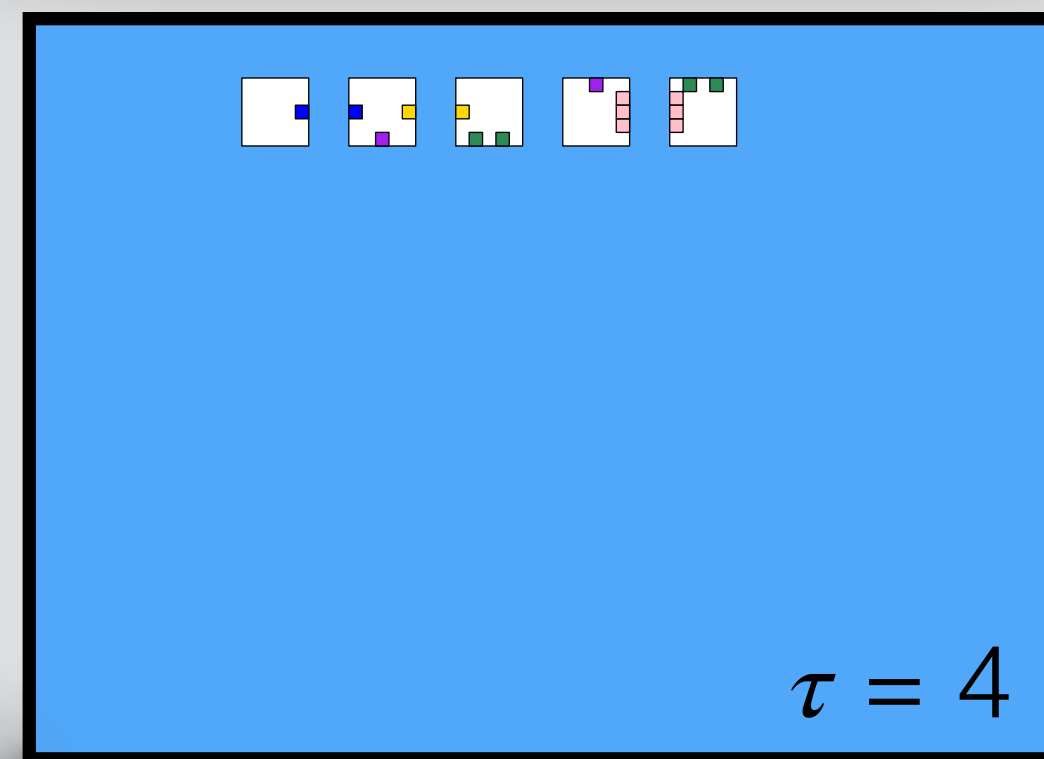
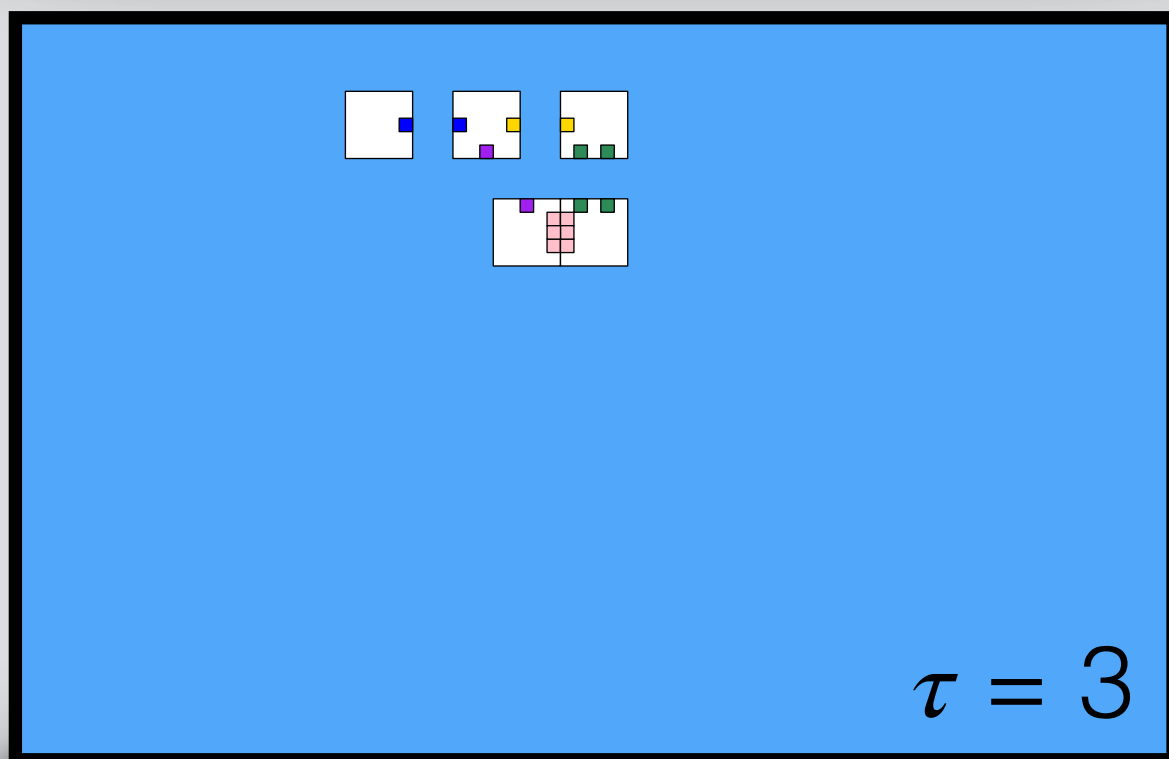
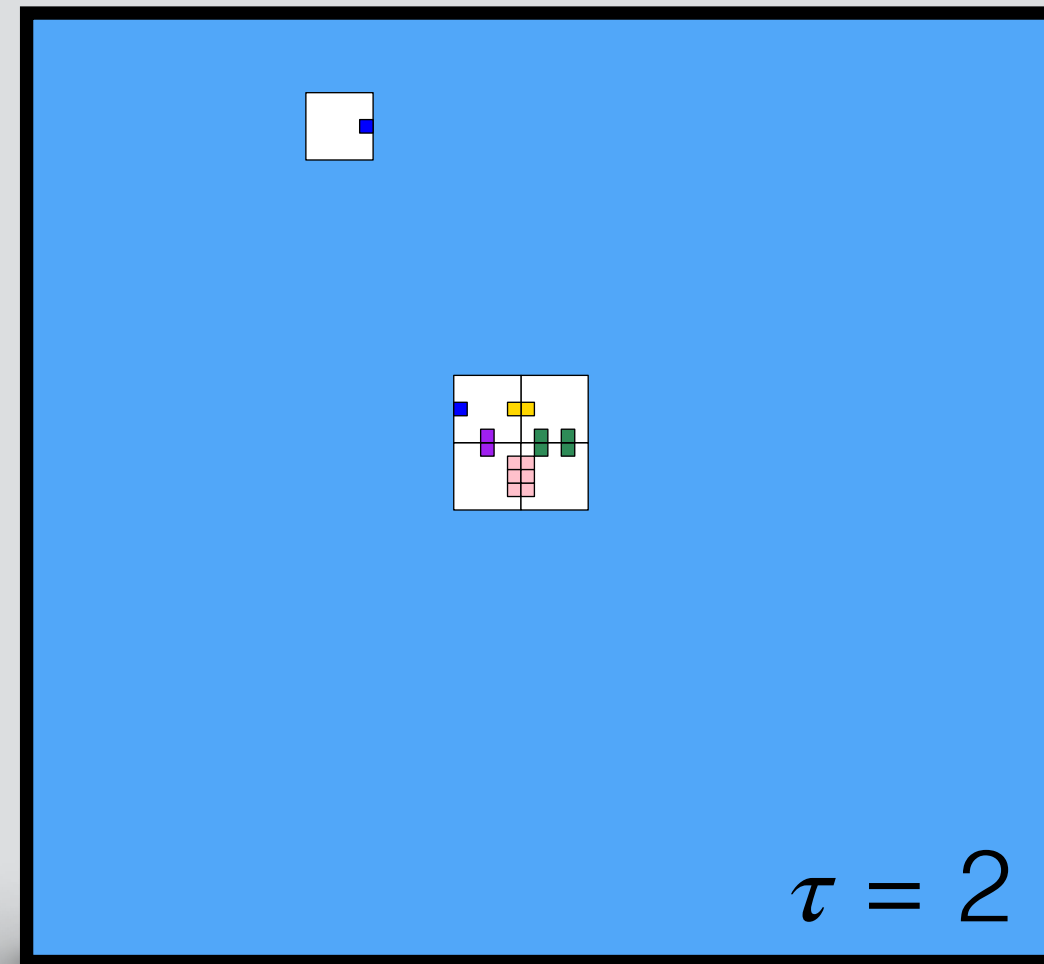
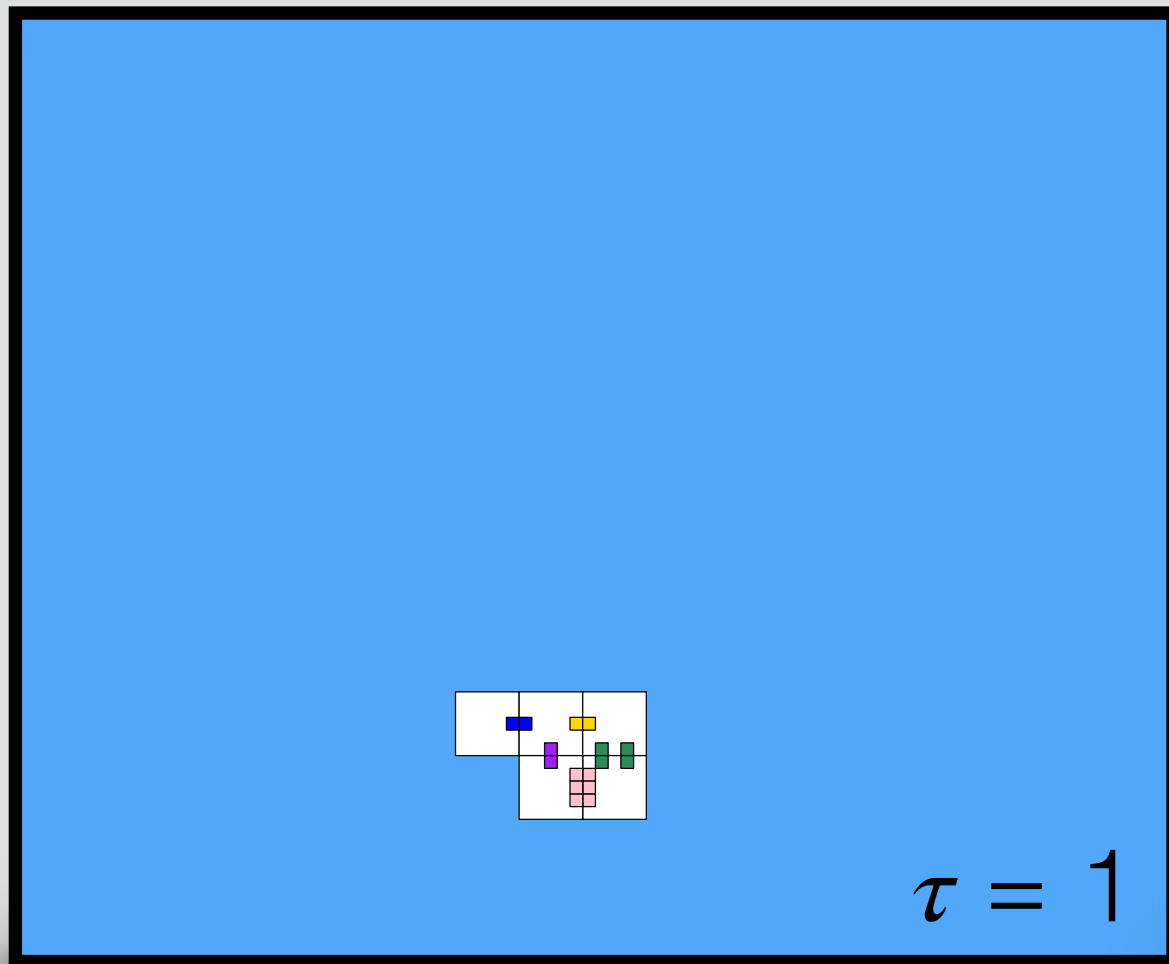
$\tau = 3$

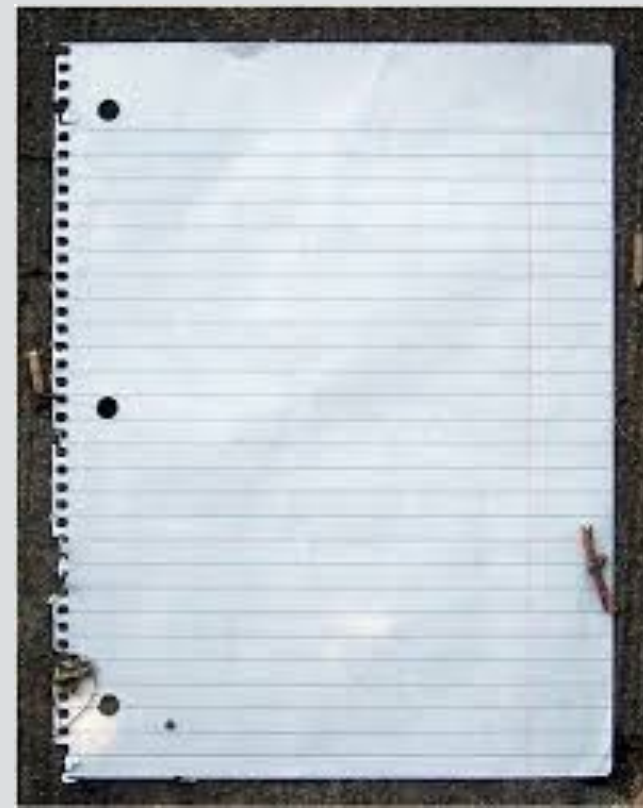
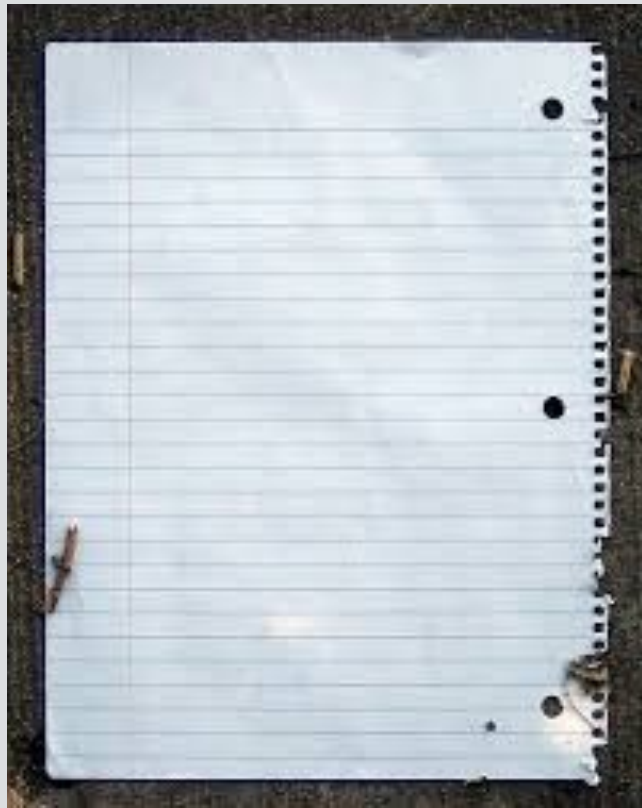


$\tau = 4$





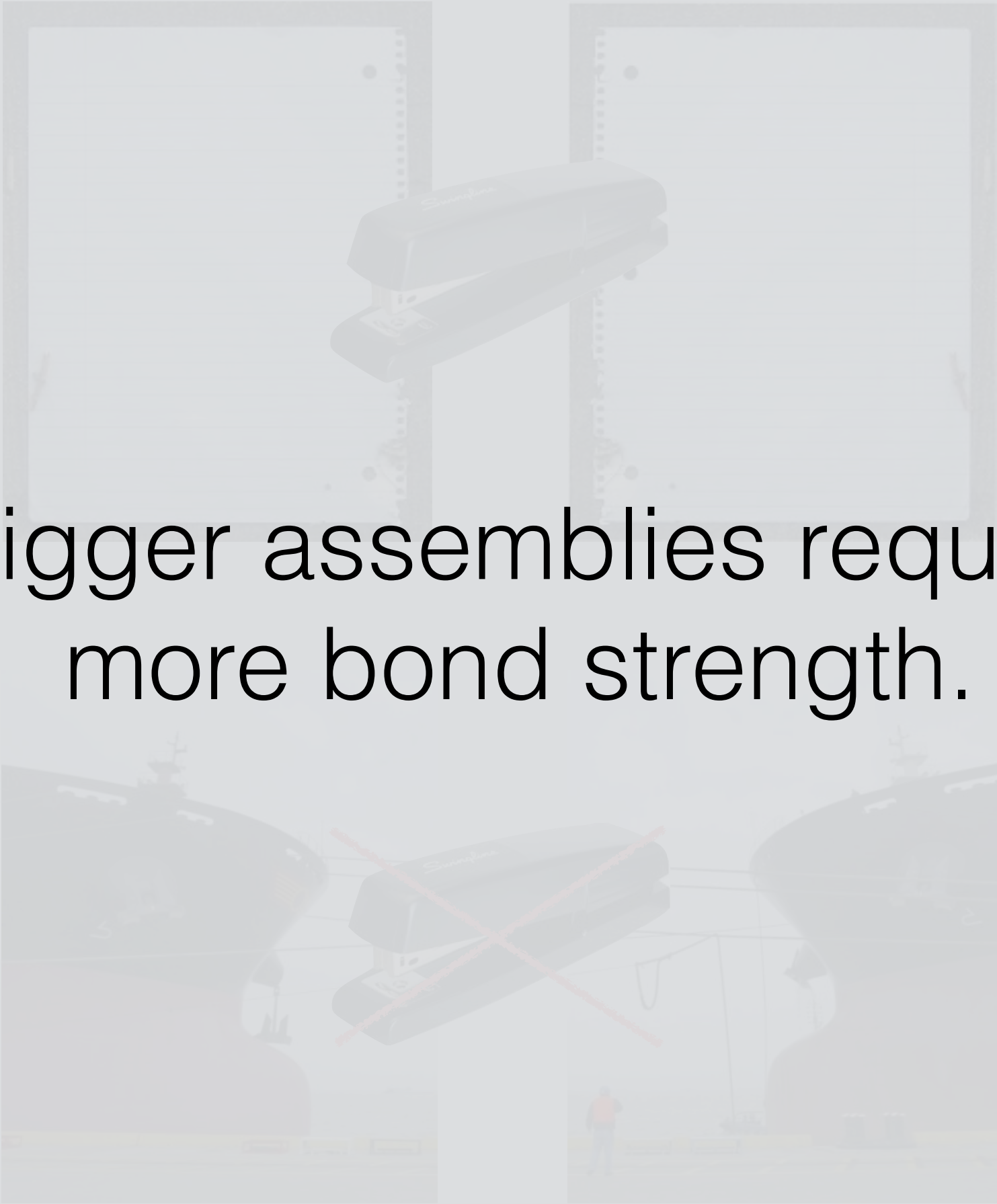












Bigger assemblies require  
more bond strength.

# Size-dependent assembly

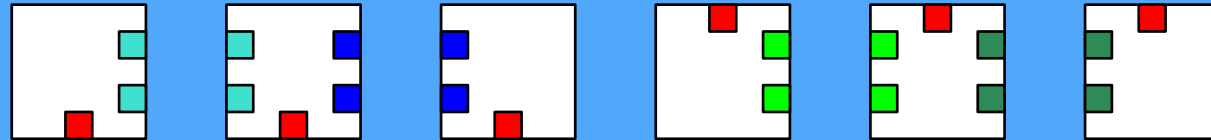
- Replace temperature  $\tau$  with increasing temperature function  $\tau : \mathbb{N} \rightarrow \mathbb{N}$ .
- Assemblies  $\alpha, \beta$  can bond if total bond strength is  $\geq \tau(\min(|\alpha|, |\beta|))$ .

# Size-dependent assembly

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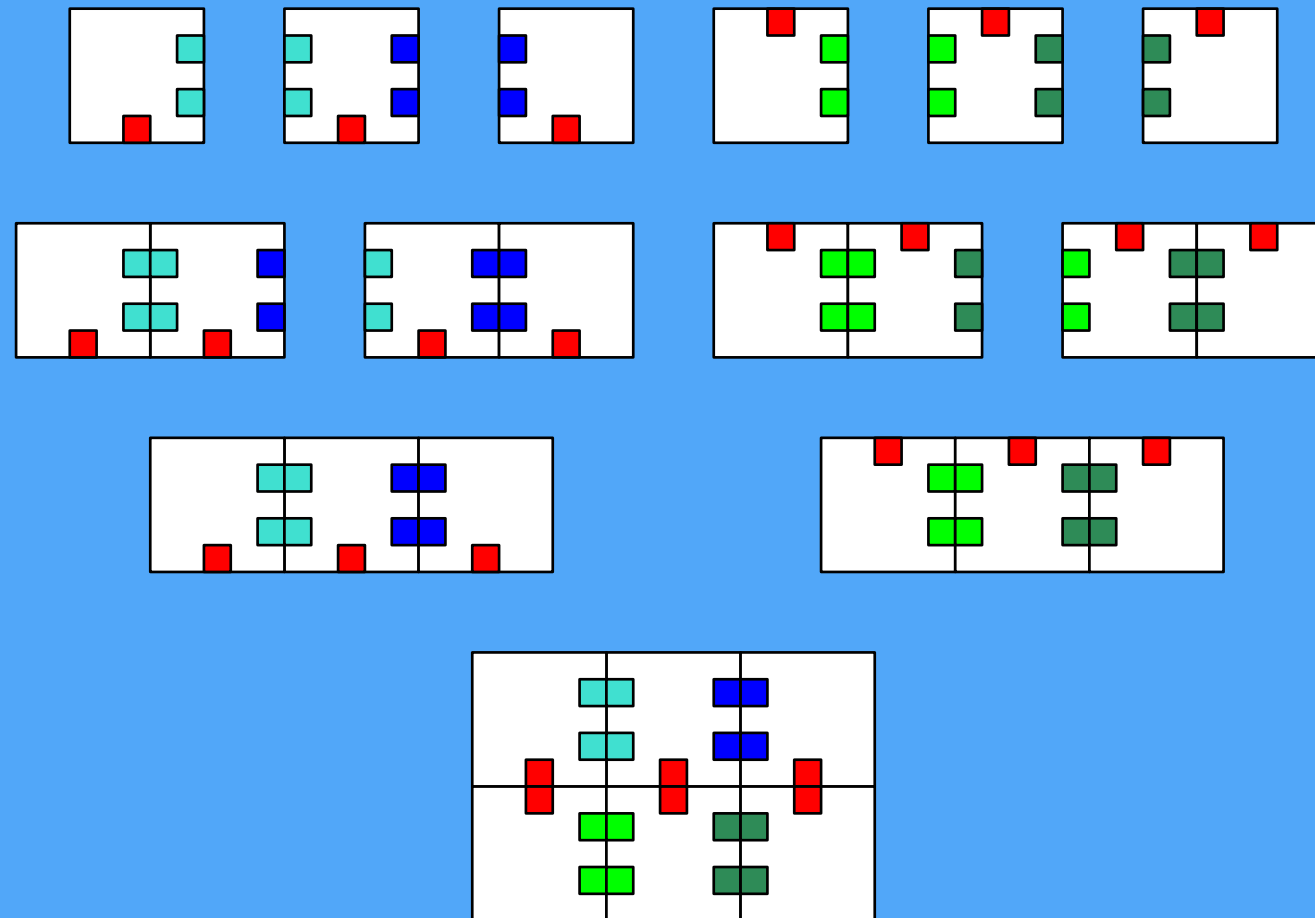


# Size-Dependent Assembly



$$\tau(n) = \begin{cases} 2 : n \leq 1 \\ 3 : \text{otherwise} \end{cases}$$

# Size-Dependent Assembly

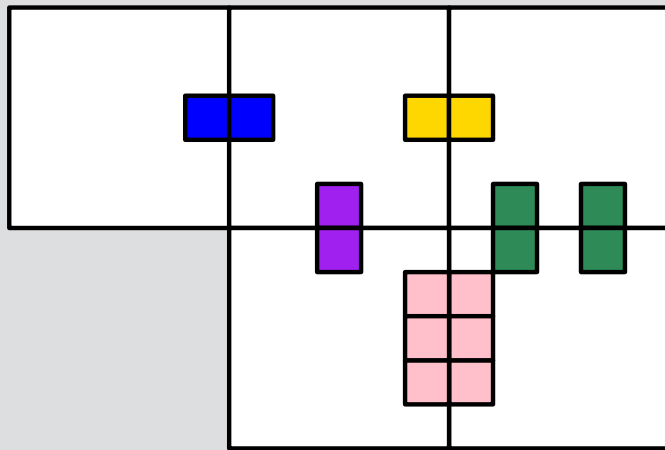


$$\tau(n) = \begin{cases} 2 & : n \leq 1 \\ 3 & : \text{otherwise} \end{cases}$$



# Stability and Cuts

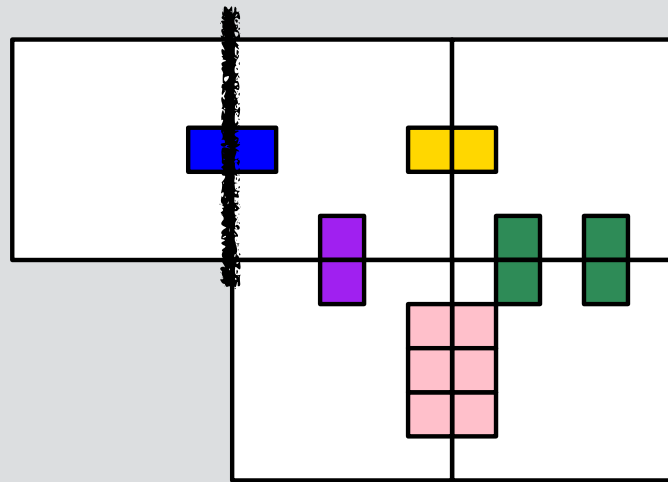
An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .



Cuts of strength:

# Stability and Cuts

An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .

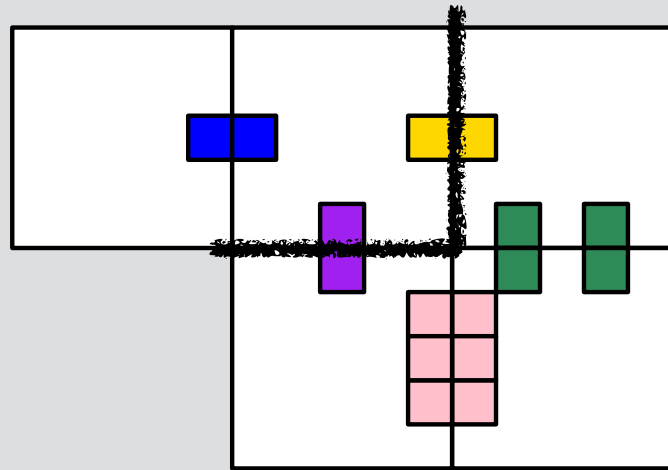


Cuts of strength:

1

# Stability and Cuts

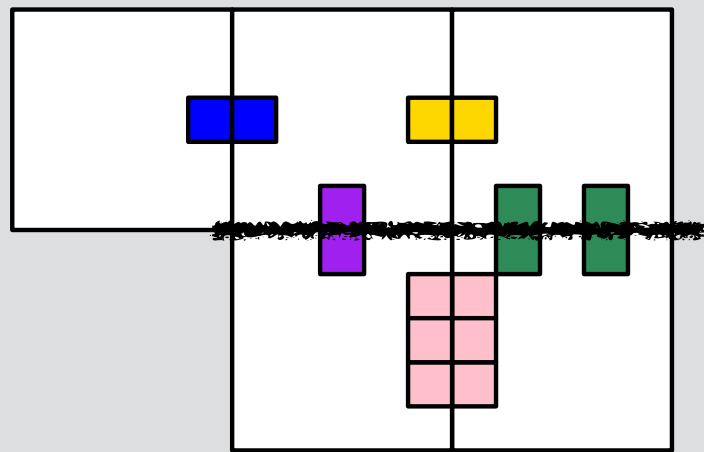
An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .



Cuts of strength:  
1, 2

# Stability and Cuts

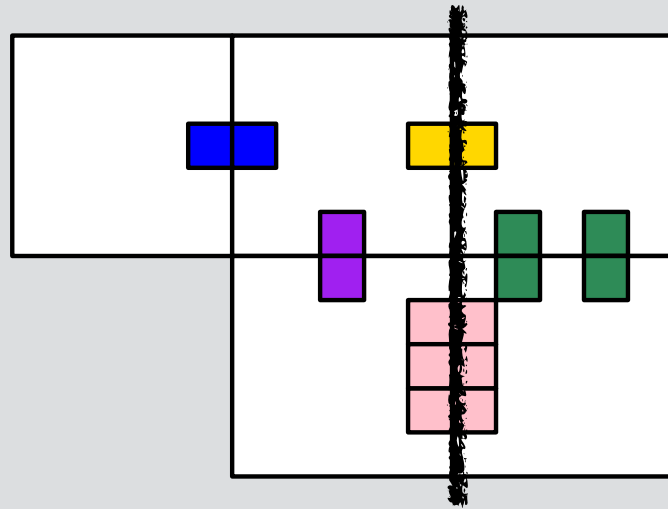
An assembly is stable at temperature  $\tau$   
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Cuts of strength:  
1, 2, 3

# Stability and Cuts

An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .



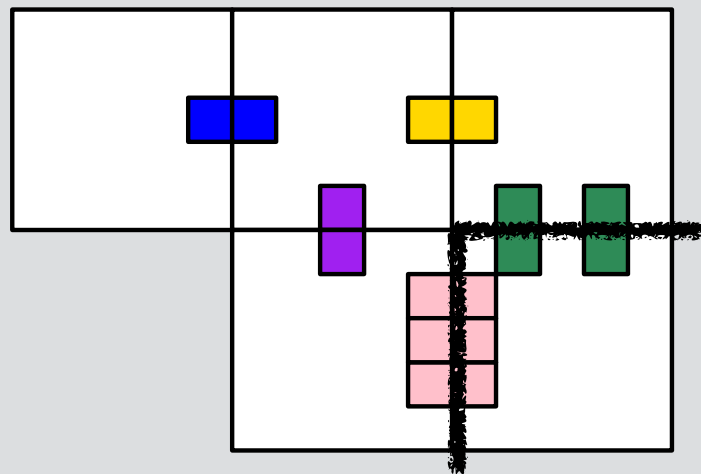
Cuts of strength:

1, 2, 3, 4



# Stability and Cuts

An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .

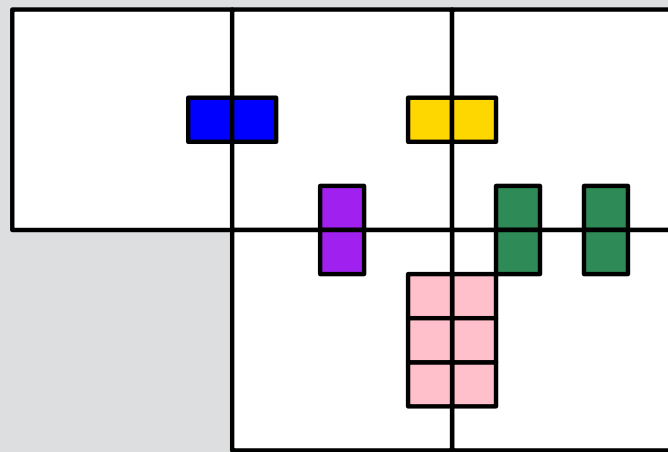


Cuts of strength:

1, 2, 3, 4, 5

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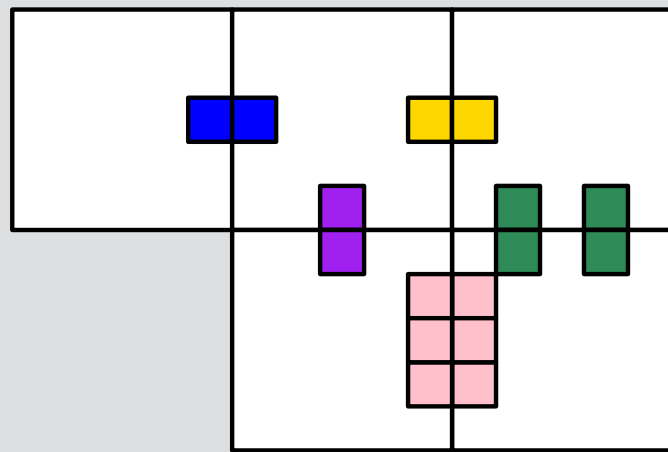


Cuts of strength:

1, 2, 3, 4, 5

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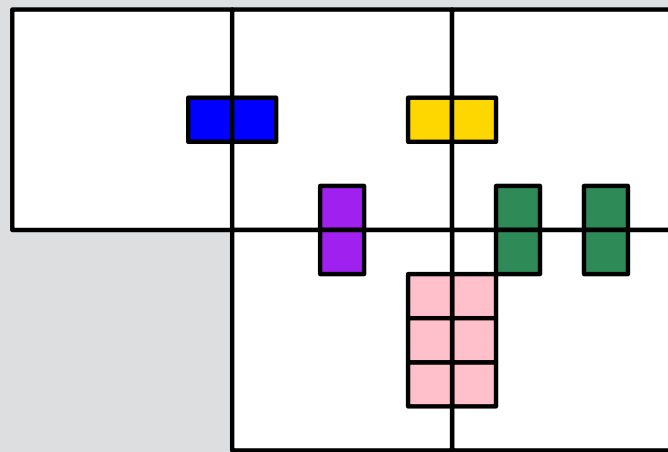
Cuts of strength:

1, 2, 3, 4, 5

Stable at  $\tau \leq 1$

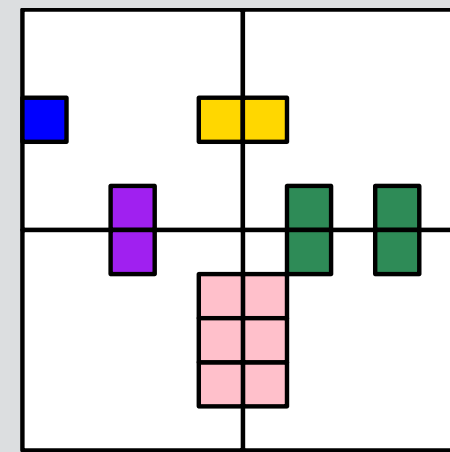
# Stability and Cuts

An assembly is stable at temperature  $\tau$   
if all cuts have strength  $\geq \tau$ .



Cuts of strength:  
1, 2, 3, 4, 5

Stable at  $\tau \leq 1$



Cuts of strength:  
2, 3, 4, 5

Stable at  $\tau \leq 2$

# Size-dependent assembly

- Replace temperature  $\tau$  with increasing temperature function  $\tau : \mathbb{N} \rightarrow \mathbb{N}$ .
- Assemblies  $\alpha, \beta$  can bond if total bond strength is  $\geq \tau(\min(|\alpha|, |\beta|))$ .

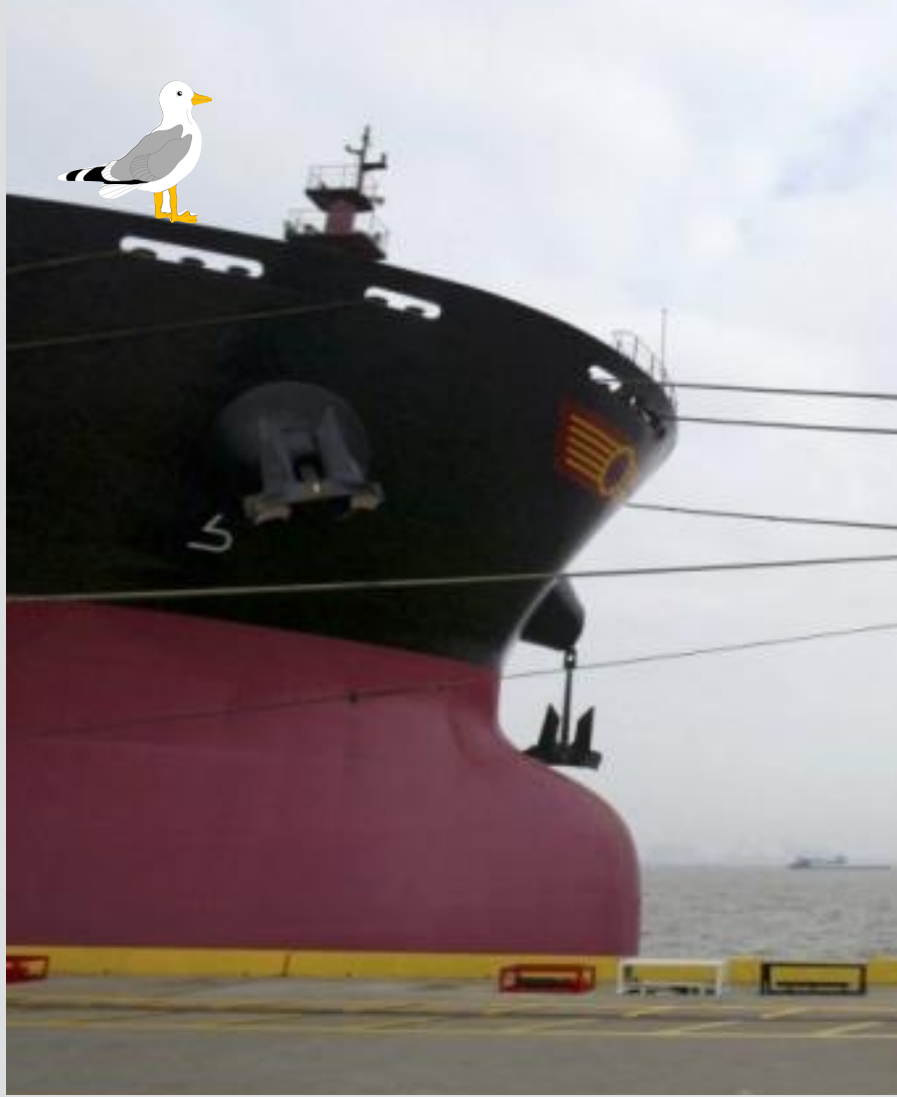
# Size-dependent assembly

- Replace temperature  $\tau$  with increasing temperature function  $\tau : \mathbb{N} \rightarrow \mathbb{N}$ .
- Assemblies  $\alpha, \beta$  can bond if total bond strength is  $\geq \tau(\min(|\alpha|, |\beta|))$ .
- Assembly is stable if every cut into connected subassemblies  $\alpha, \beta$  has strength  $\geq \tau(\min(|\alpha|, |\beta|))$ .

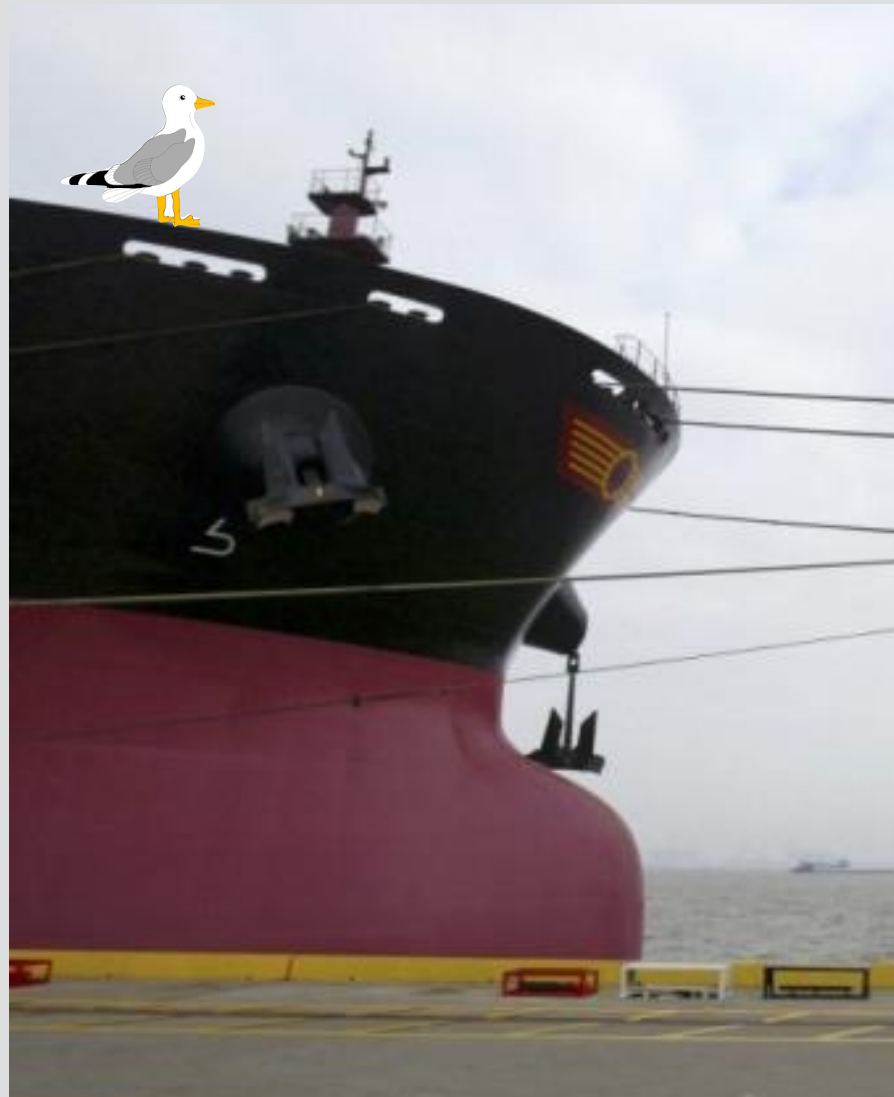




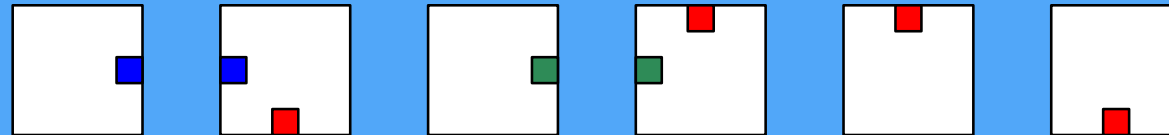




Unstable assemblies break along weak cuts.

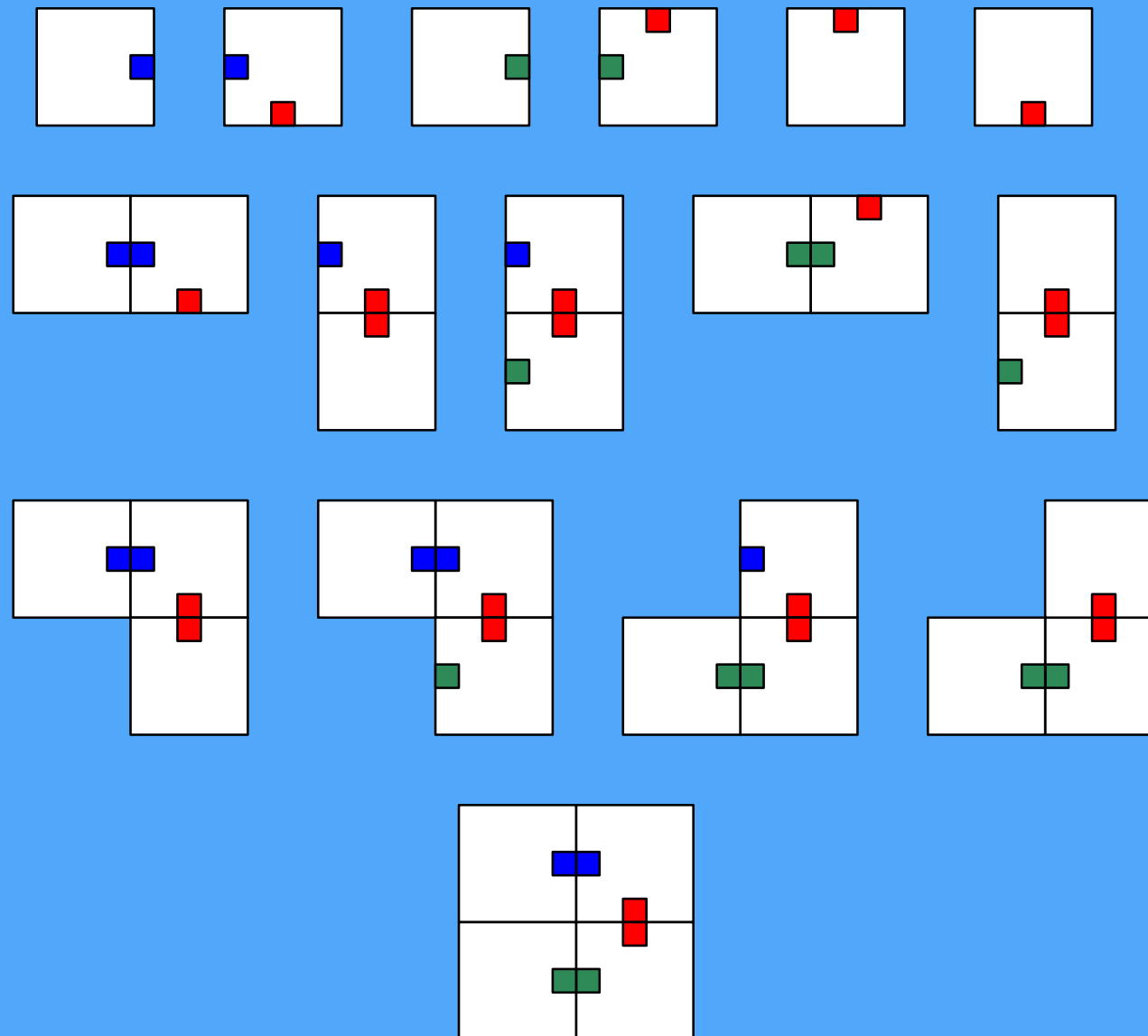


# Size-Dependent Assembly



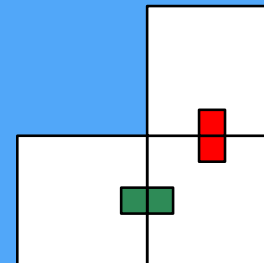
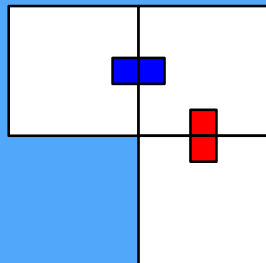
$$\tau(n) = \begin{cases} 1 : n \leq 1 \\ 2 : \text{otherwise} \end{cases}$$

# Size-Dependent Assembly



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# Size-Dependent Assembly



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# Questions

Can temperature functions do anything “useful”,  
e.g. build shapes more efficiently?

Breakage looks complicated.

How hard is deciding if an assembly is stable?

# Questions and Prior Work

Can temperature functions do anything “useful”,  
e.g. build shapes more efficiently?

For fixed  $\tau$ ,  $N \times N$ :  $\Theta(\log(N)/\log\log(N))$  tile types,  
 $C \times N$ :  $\Theta(N^{1/C})$  tile types.

Breakage looks complicated.

How hard is deciding if an assembly is stable?

For fixed  $\tau$ , polynomial-time (min-cut).

# Questions and Answers

Can temperature functions do anything “useful”,  
e.g. build shapes more efficiently?

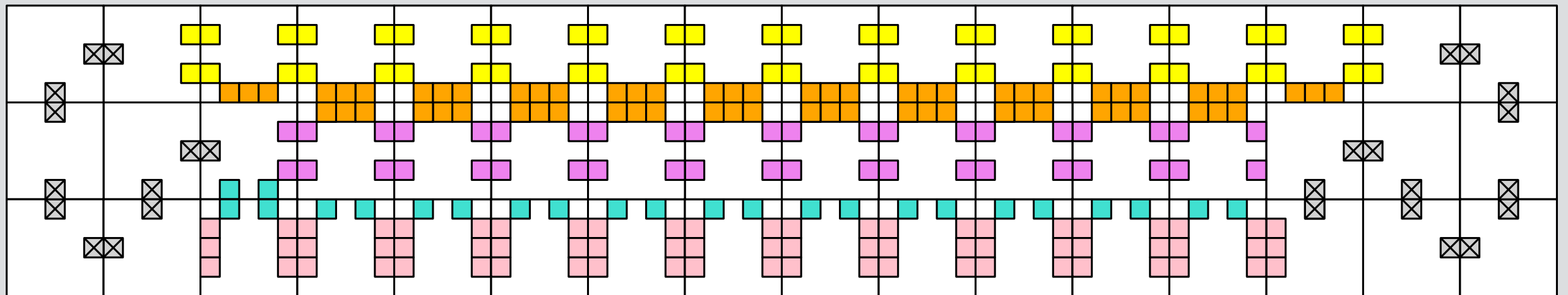
There exists a set of tiles  $T$  that assembles  $3 \times N$   
rectangle for each  $N \geq 7$ , given appropriate  $\tau(n)$ .

Breakage looks complicated.

How hard is deciding if an assembly is stable?

coNP-complete.

# 3xN Rectangle Construction



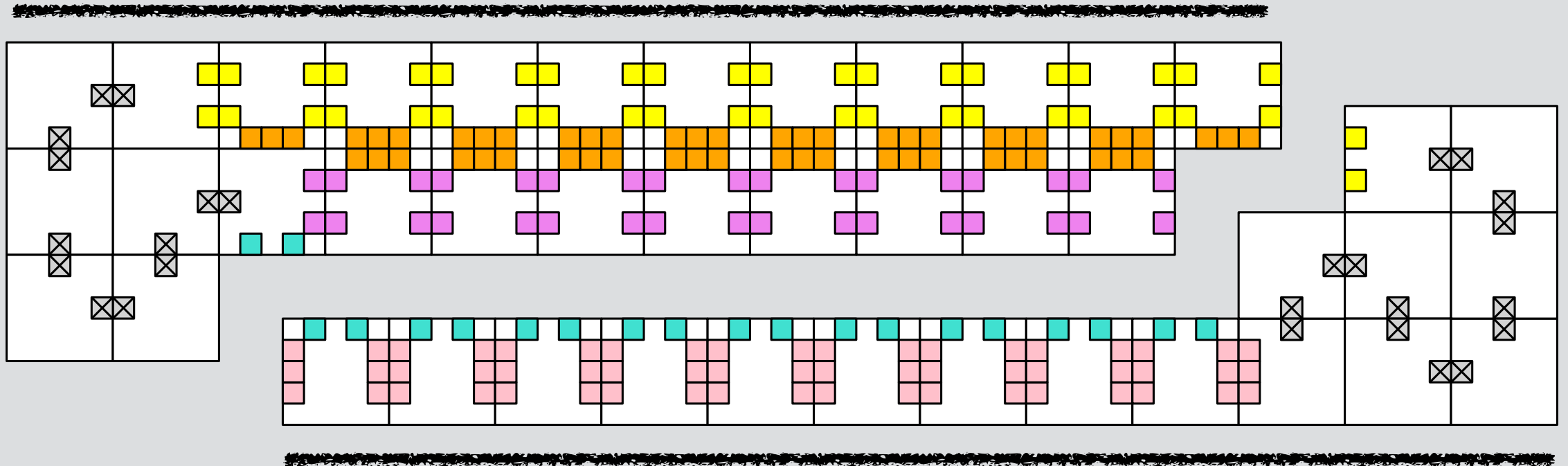
*Terminal assembly*

$$\tau(n) = \begin{cases} 3 : n \leq N - 6 \\ 4 : N - 5 \leq n \leq N + 3 \\ 5 : N + 4 \leq n \leq 2N - 2 \\ 8 : \text{otherwise} \end{cases}$$

*Temperature function*

# 3xN Rectangle Construction

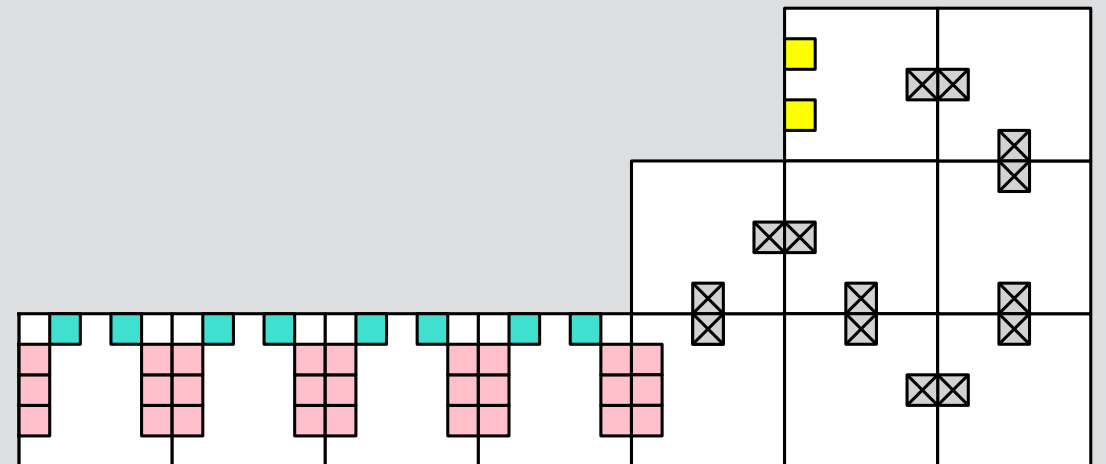
$$\geq N - 2$$



$$\leq N - 2$$

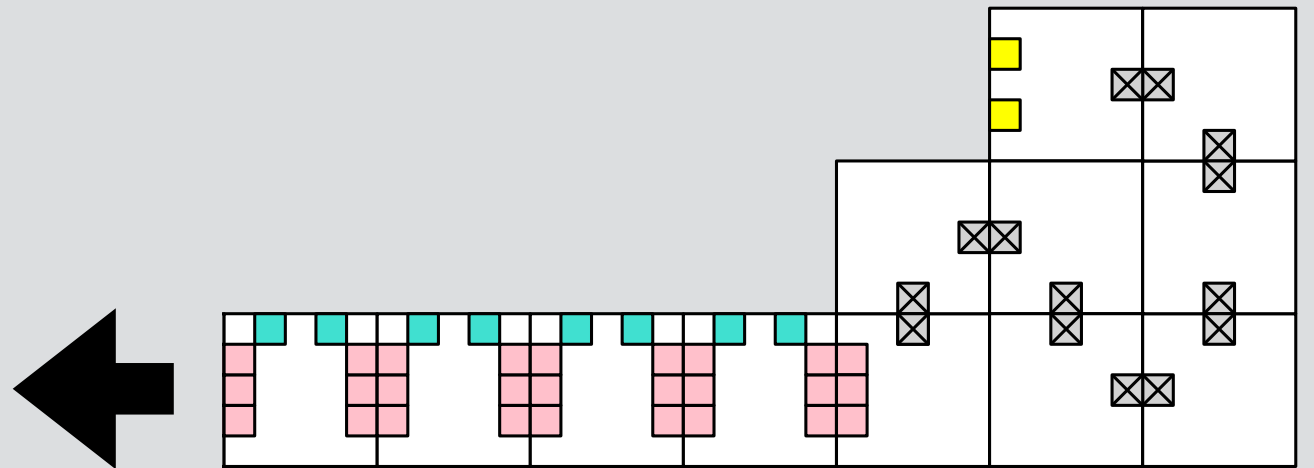
$$\tau(n) = \begin{cases} 3 : n \leq N - 6 \\ 4 : N - 5 \leq n \leq N + 3 \\ 5 : N + 4 \leq n \leq 2N - 2 \\ 8 : \text{otherwise} \end{cases}$$

# 3xN Rectangle Construction



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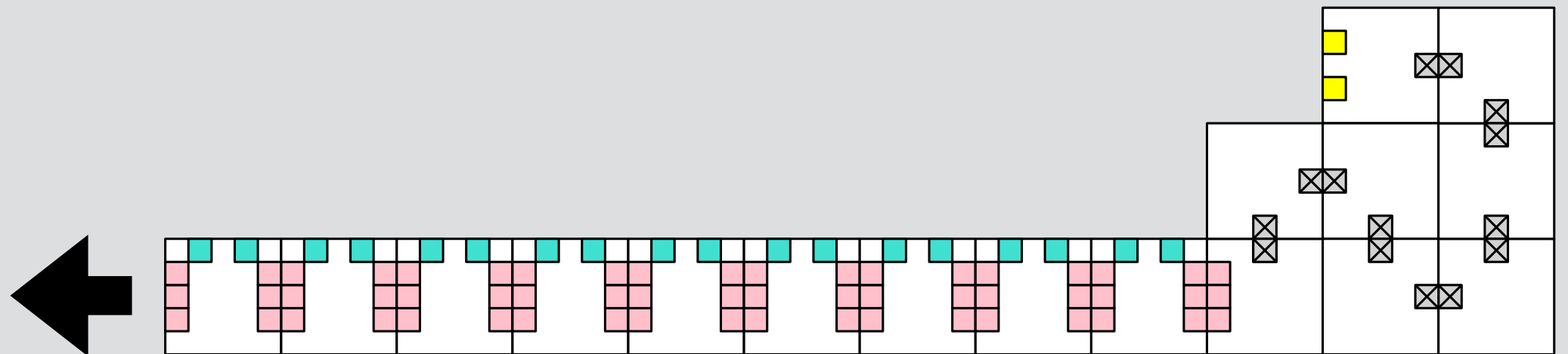
# 3xN Rectangle Construction



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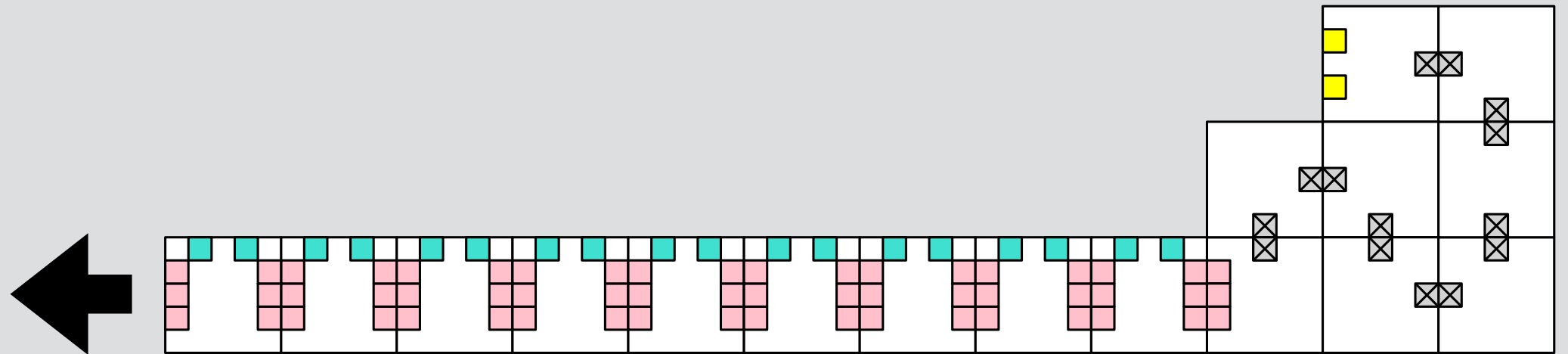


# 3xN Rectangle Construction



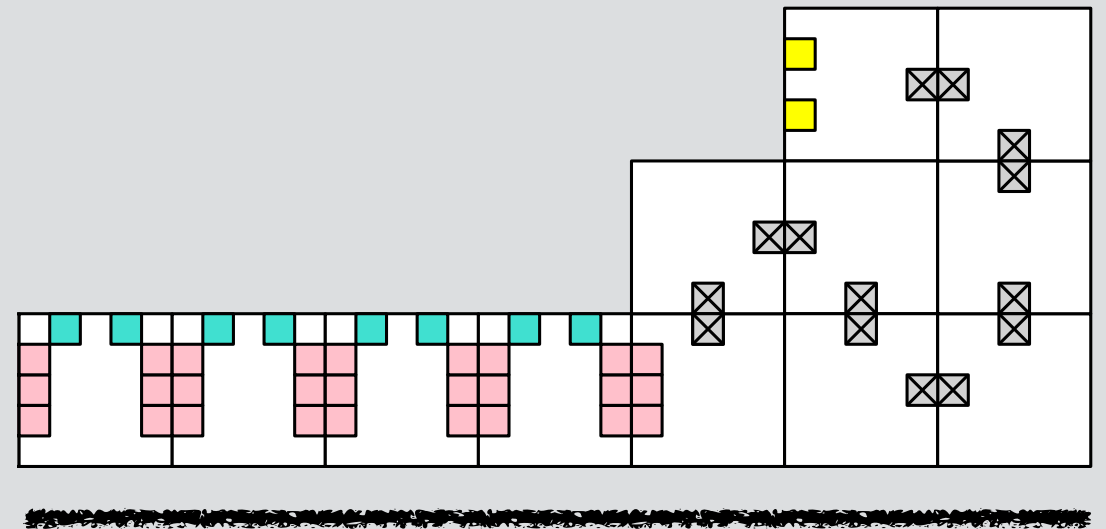
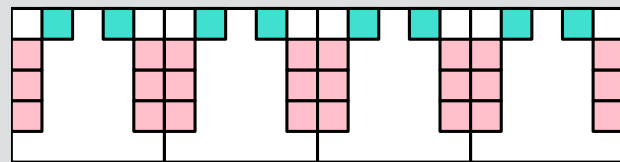
$$\tau(n) = \begin{cases} 3 : n \leq N - 6 \\ 4 : N - 5 \leq n \leq N + 3 \\ 5 : N + 4 \leq n \leq 2N - 2 \\ 8 : \text{otherwise} \end{cases}$$

# 3xN Rectangle Construction



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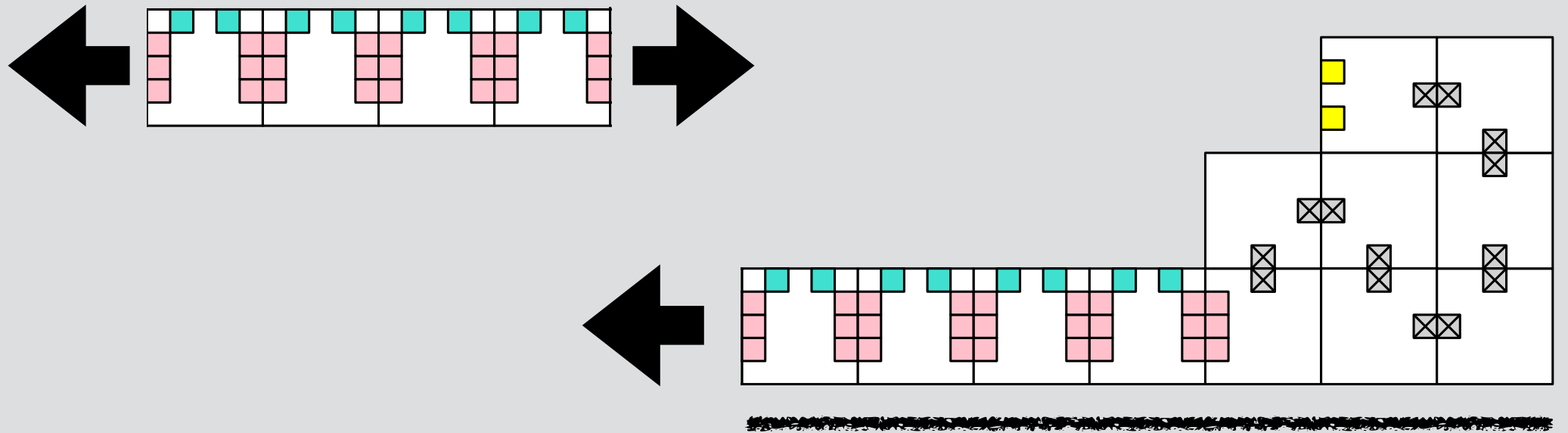
# 3xN Rectangle Construction



$$\leq N - 2$$

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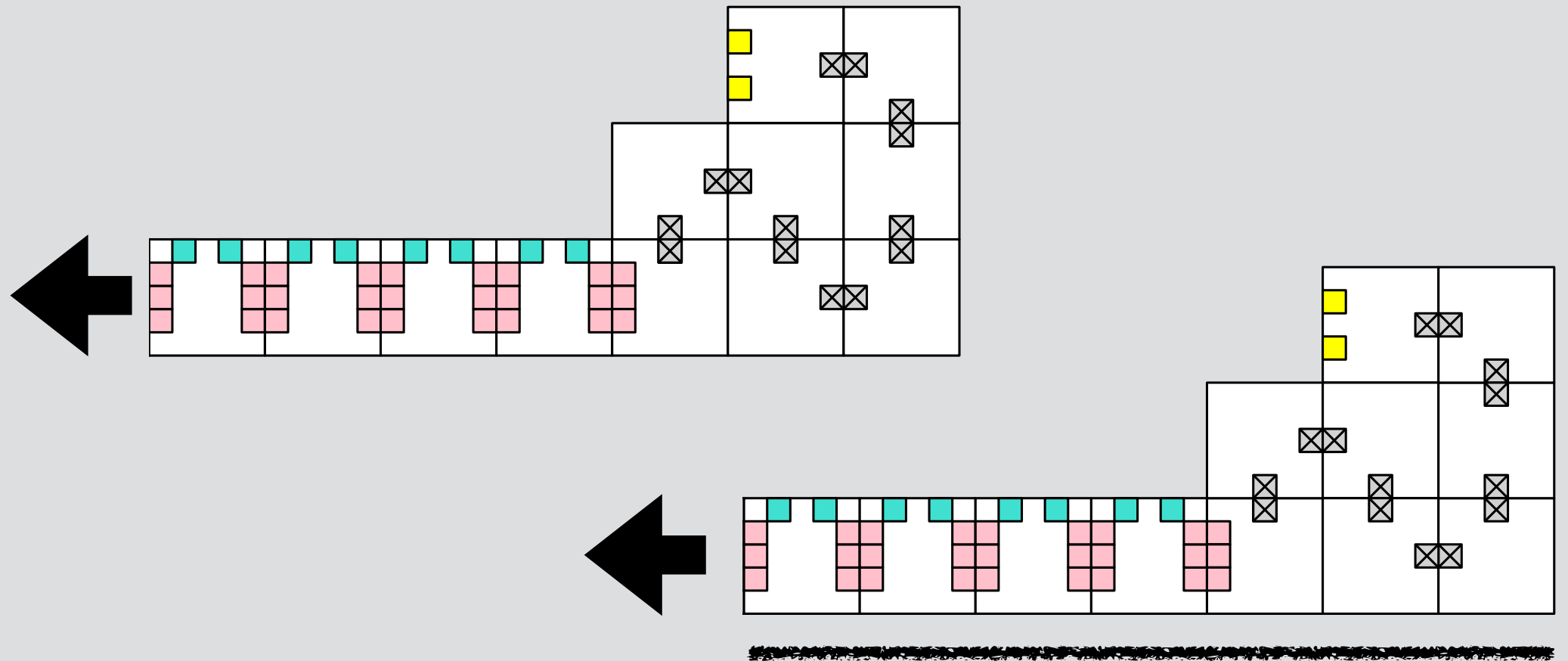
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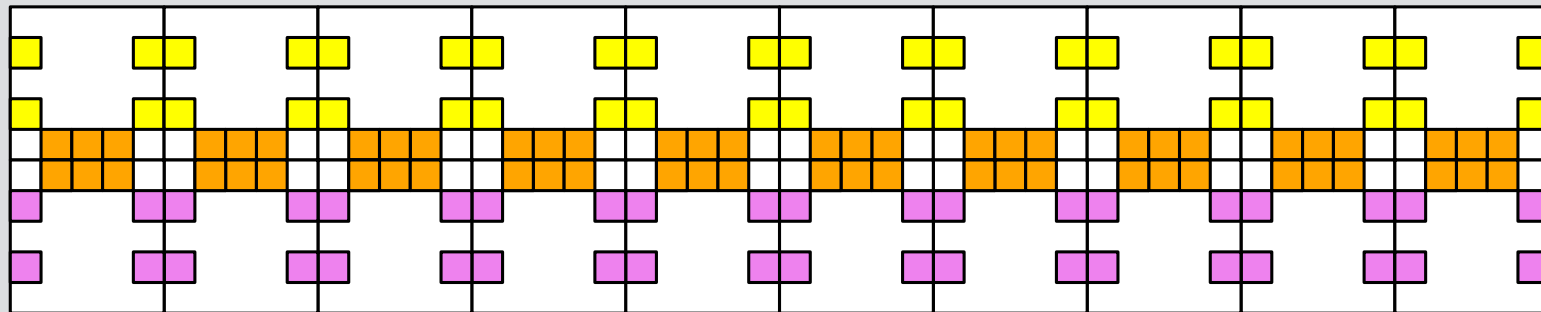
# 3xN Rectangle Construction



$\leq N - 2$

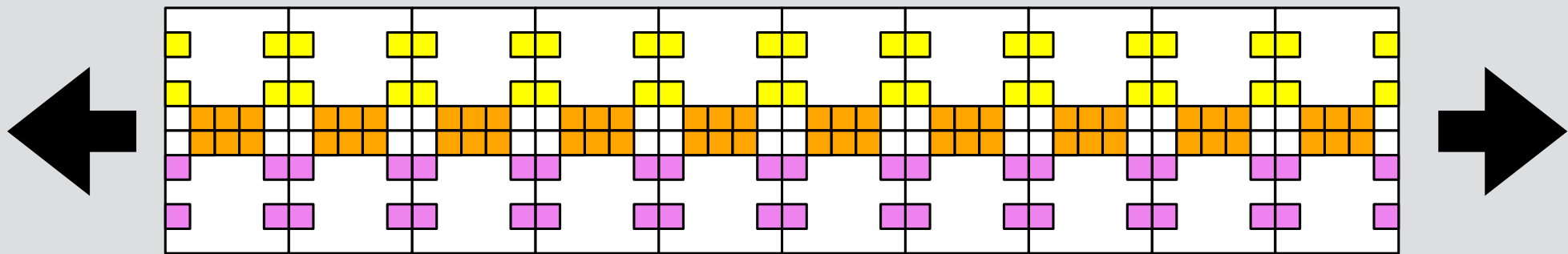
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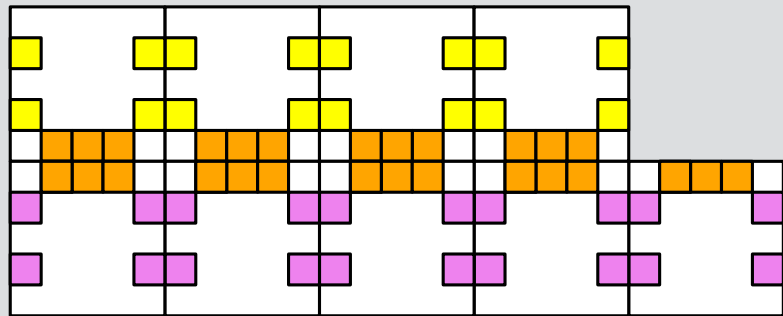
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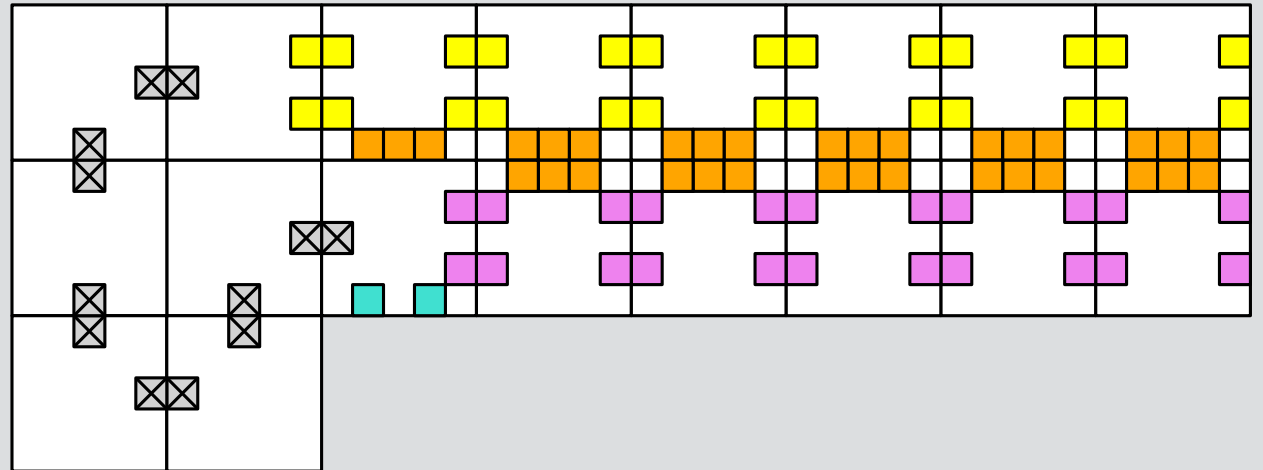
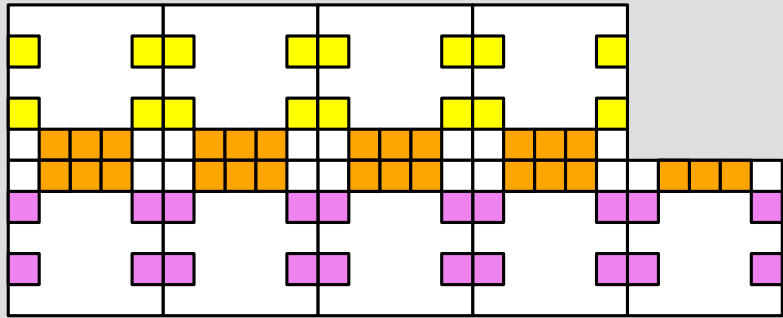
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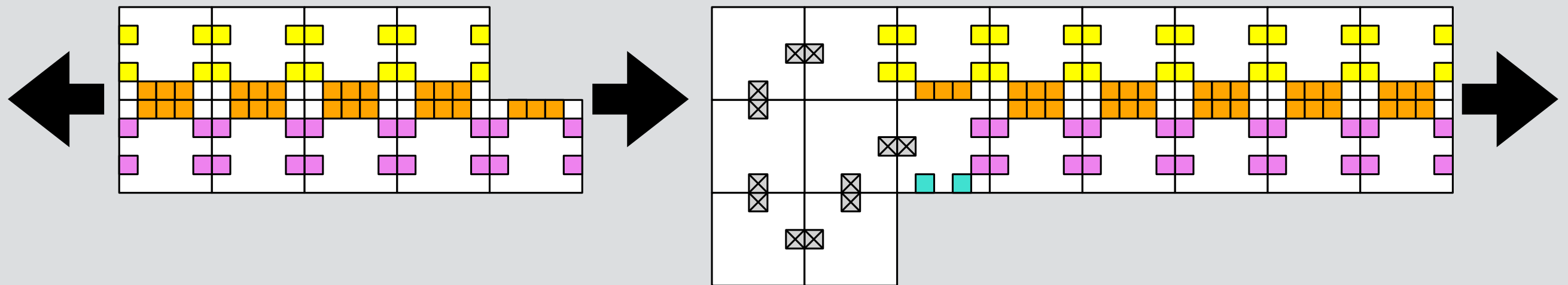


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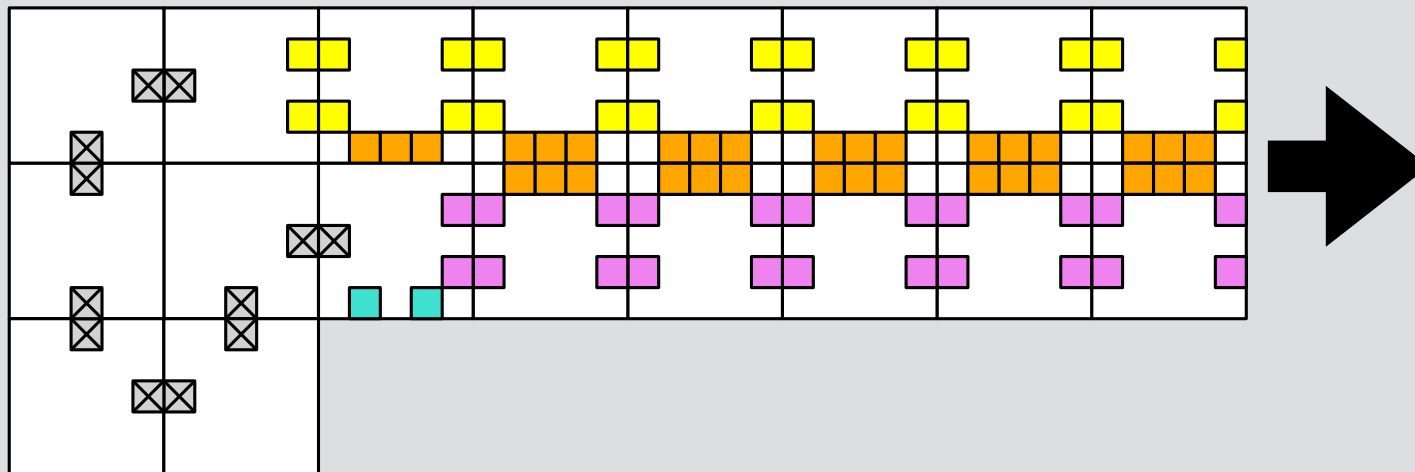
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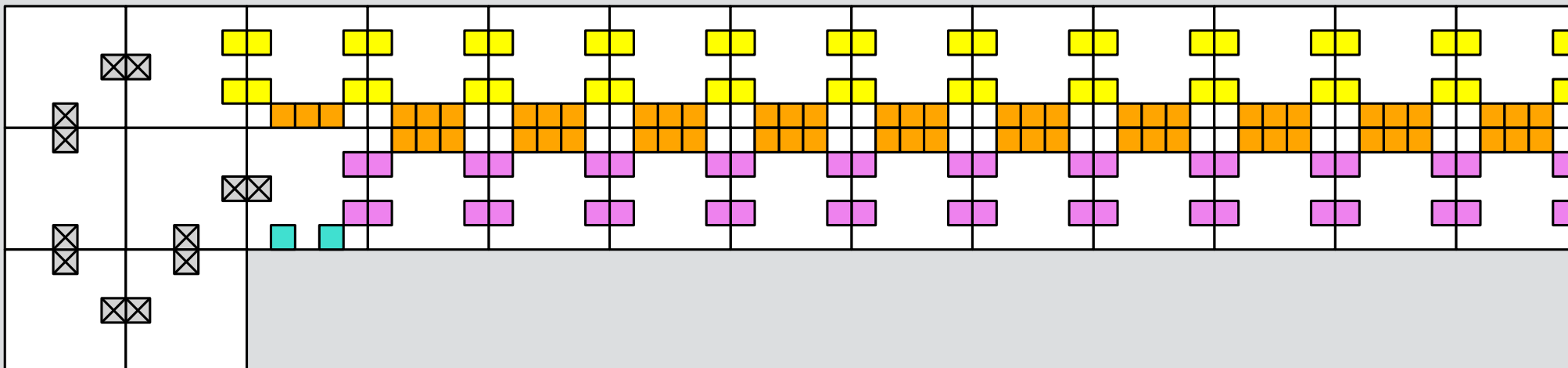
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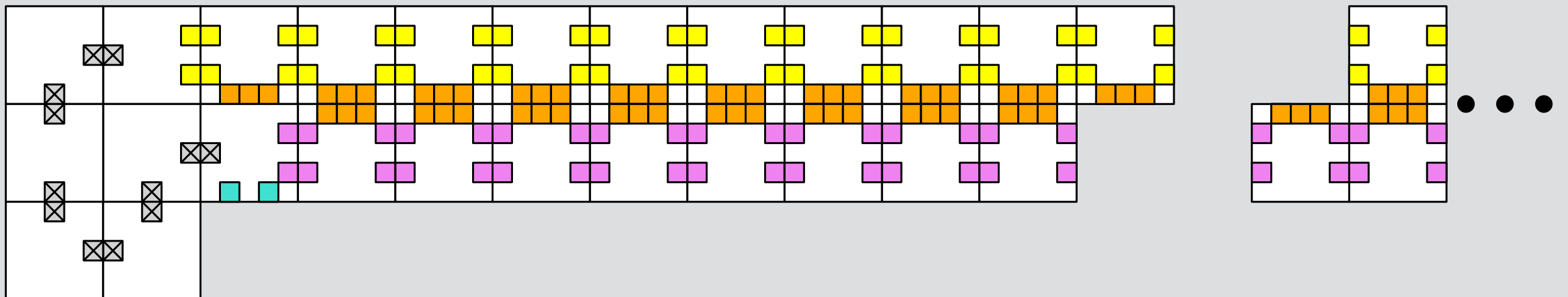
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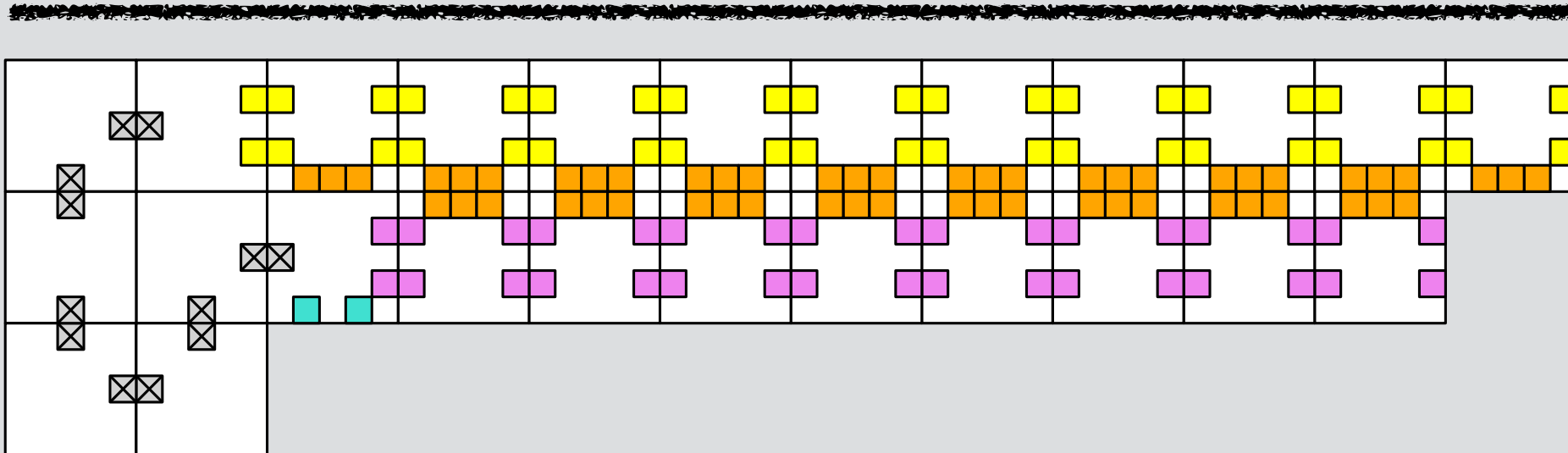
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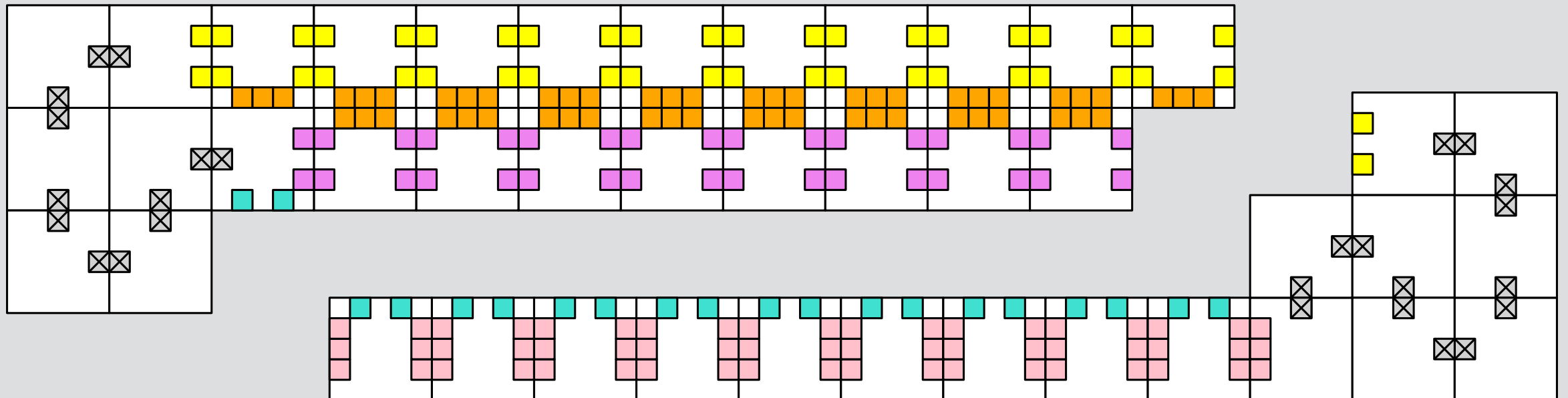
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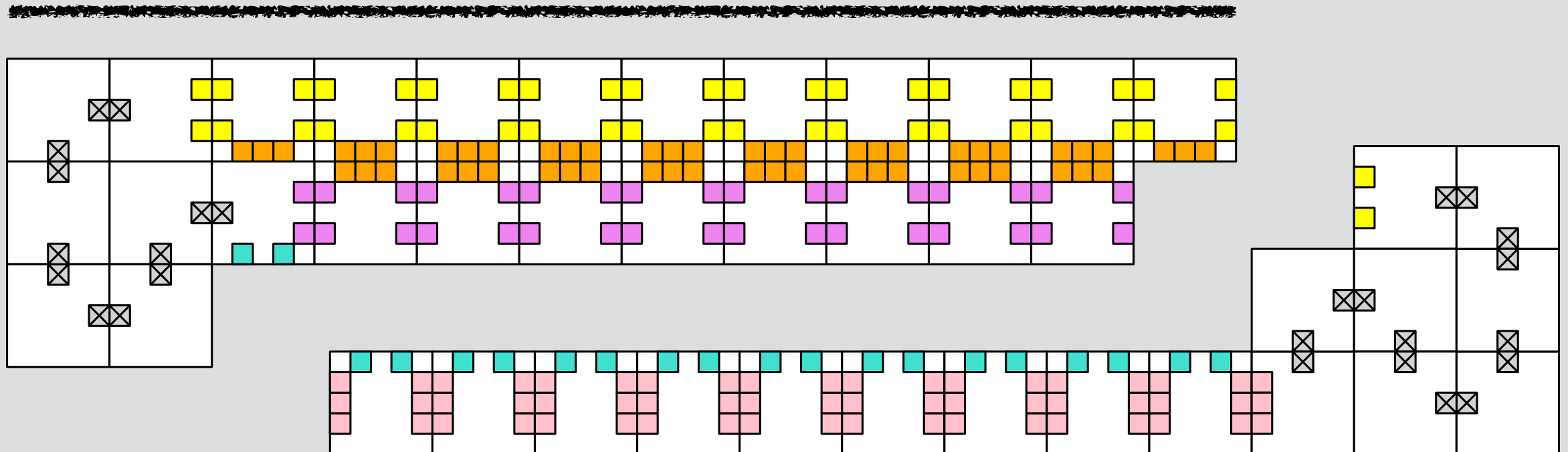
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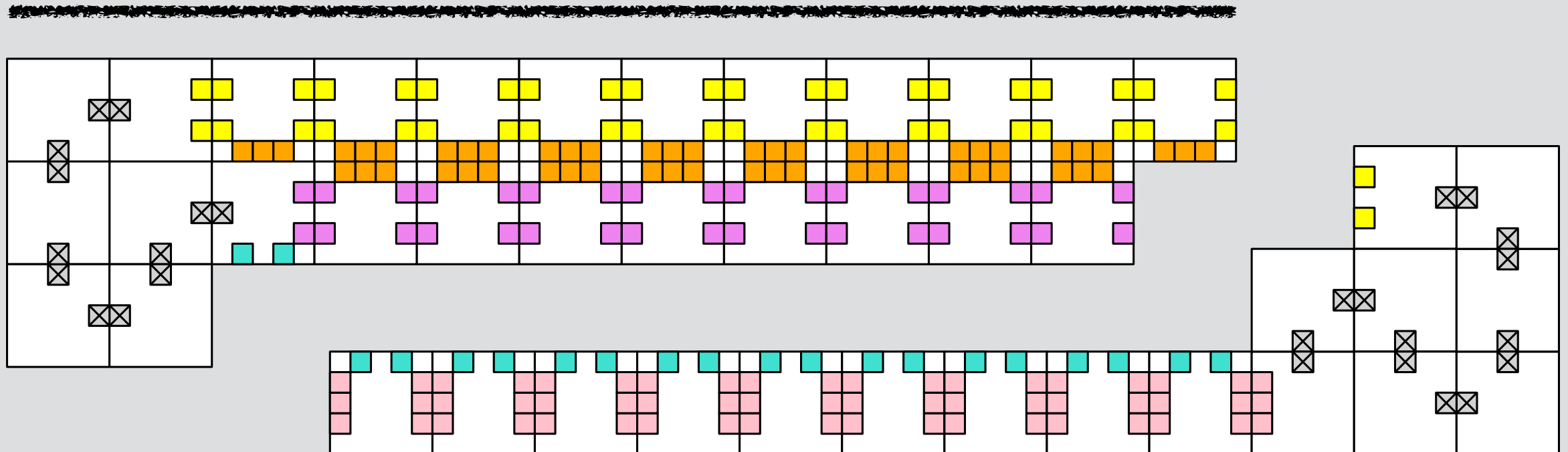
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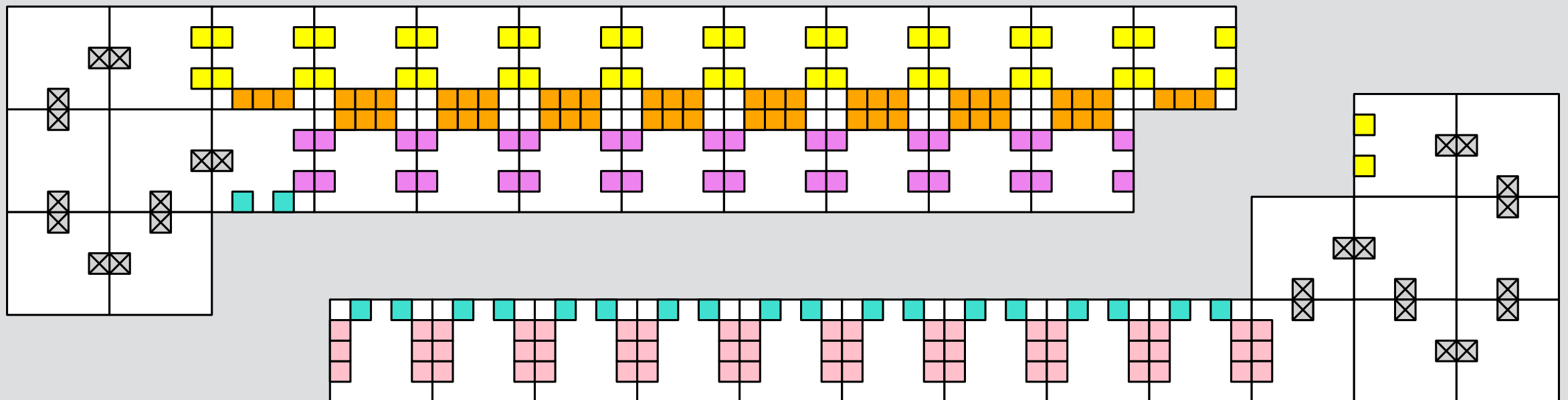
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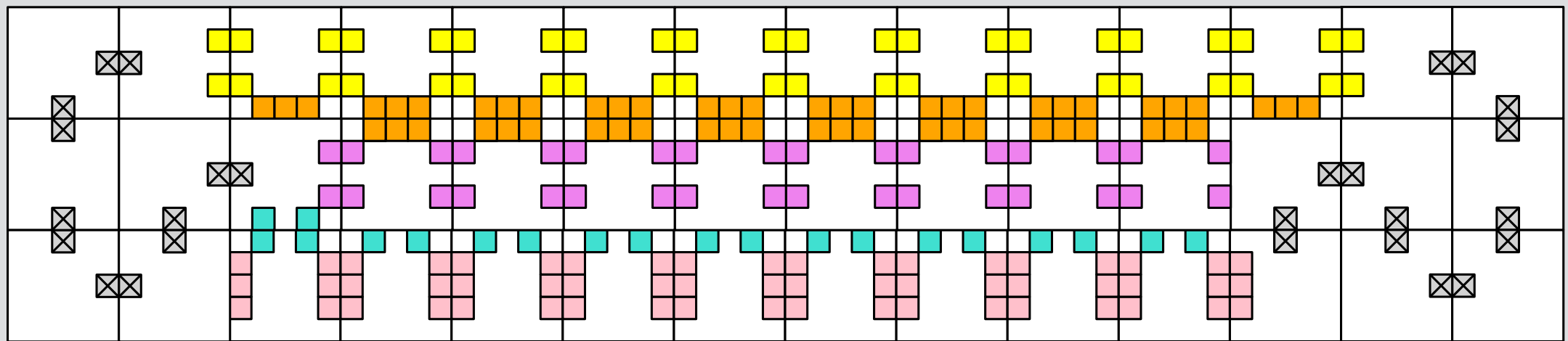


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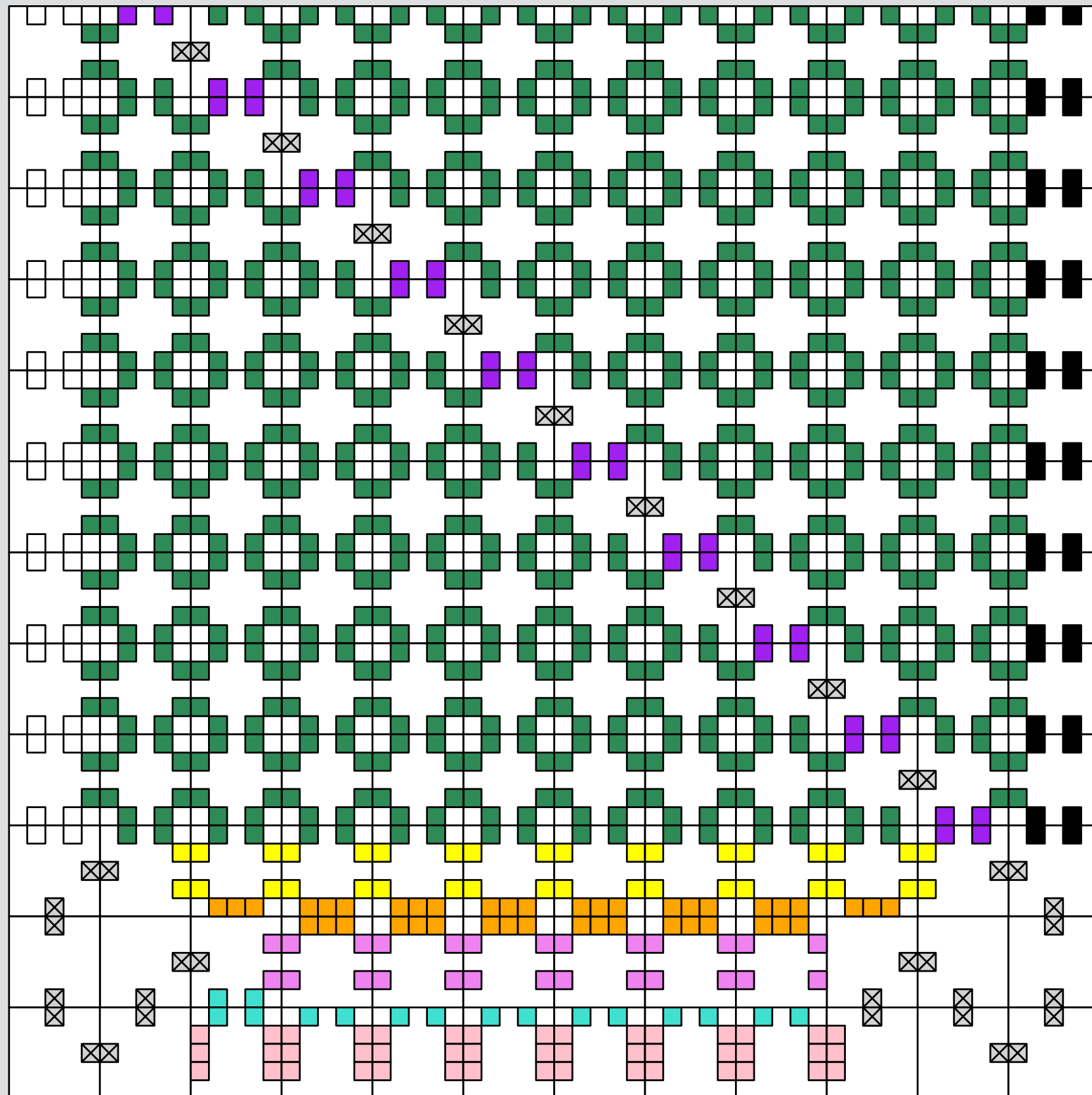
# 3xN Rectangle Construction

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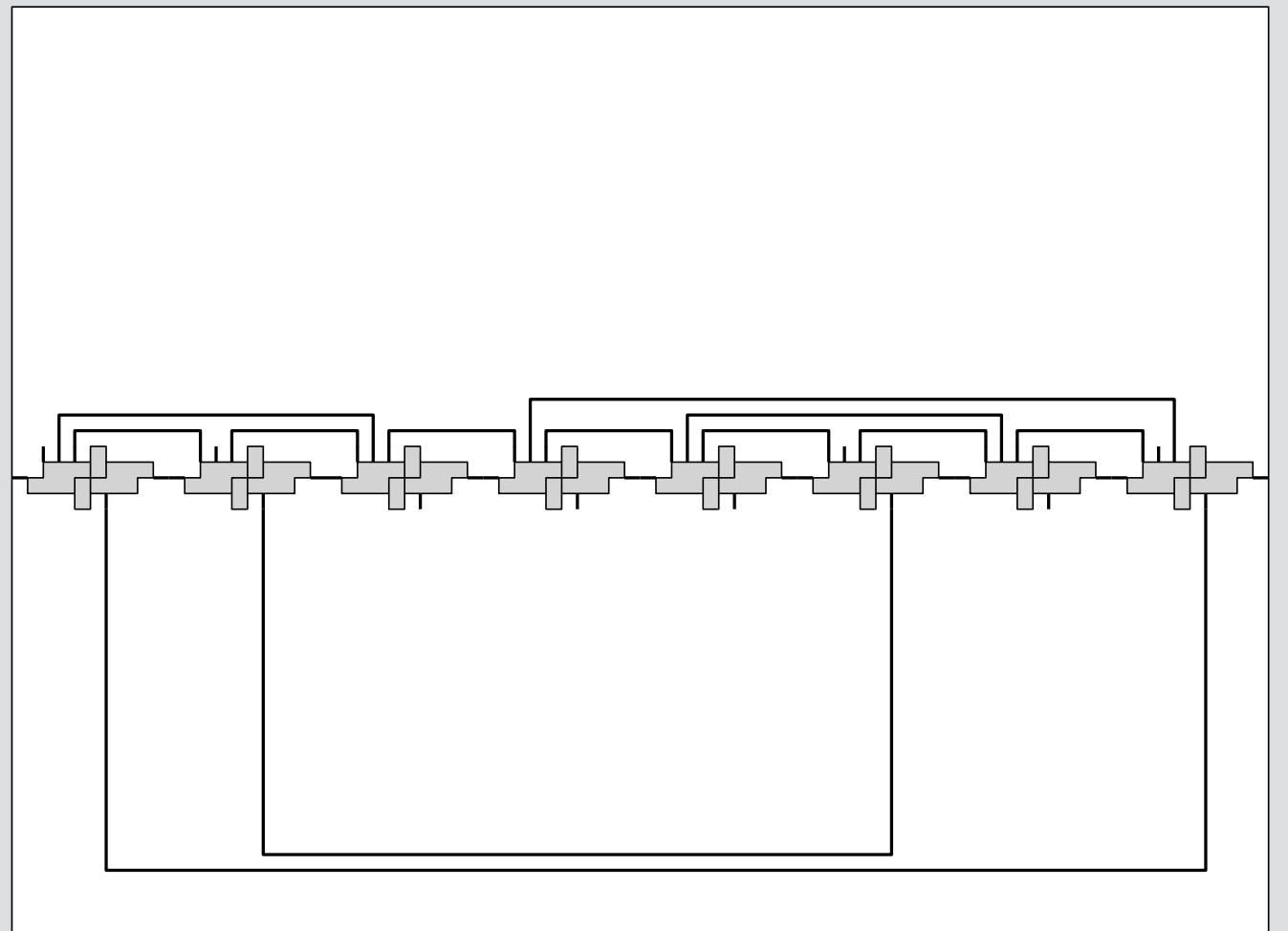
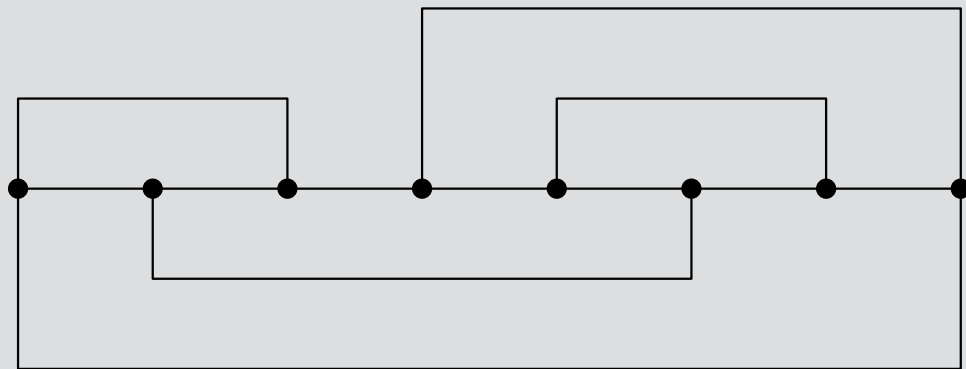
# NxN Square Construction



Same  $\tau(n)$

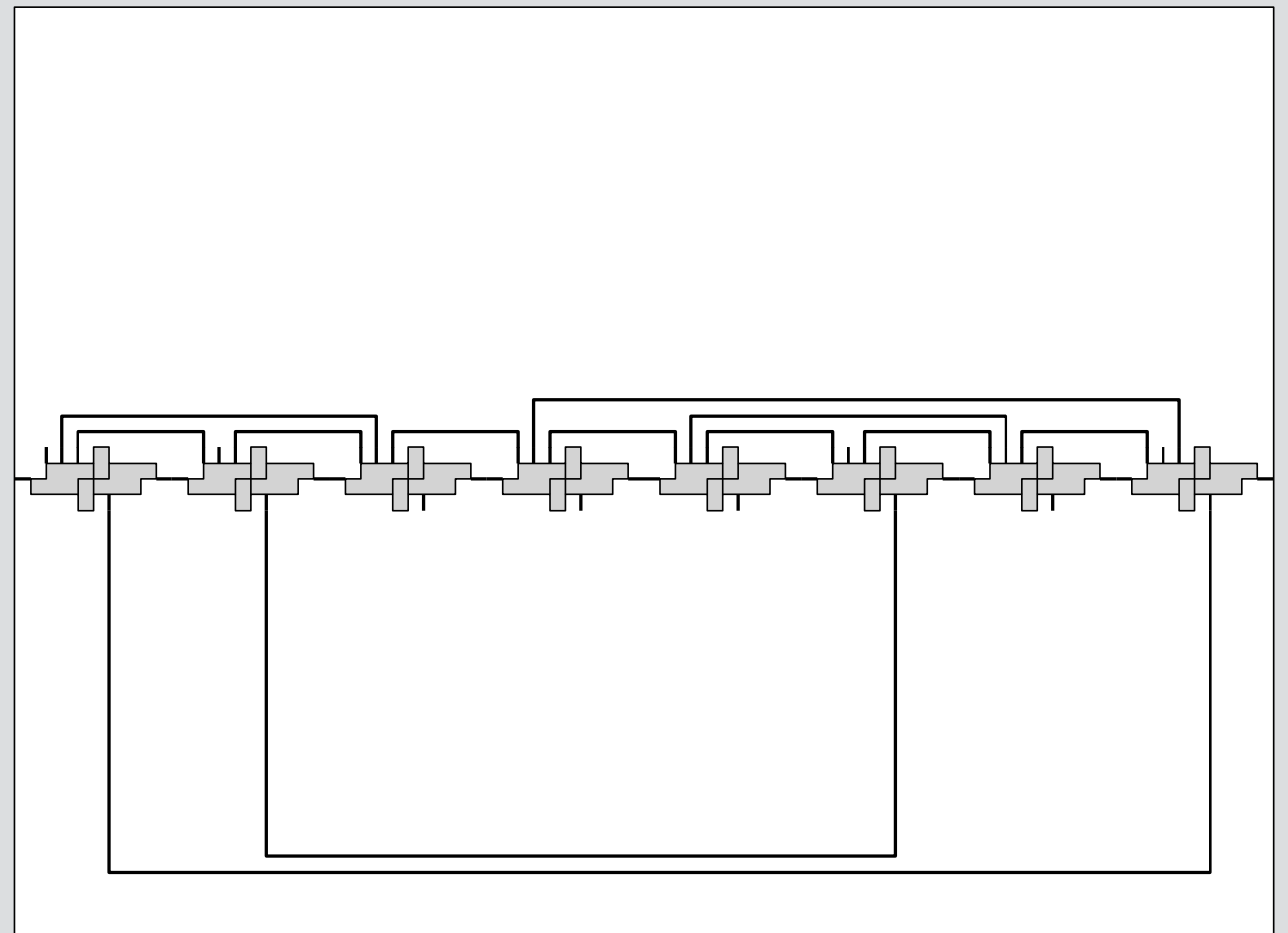
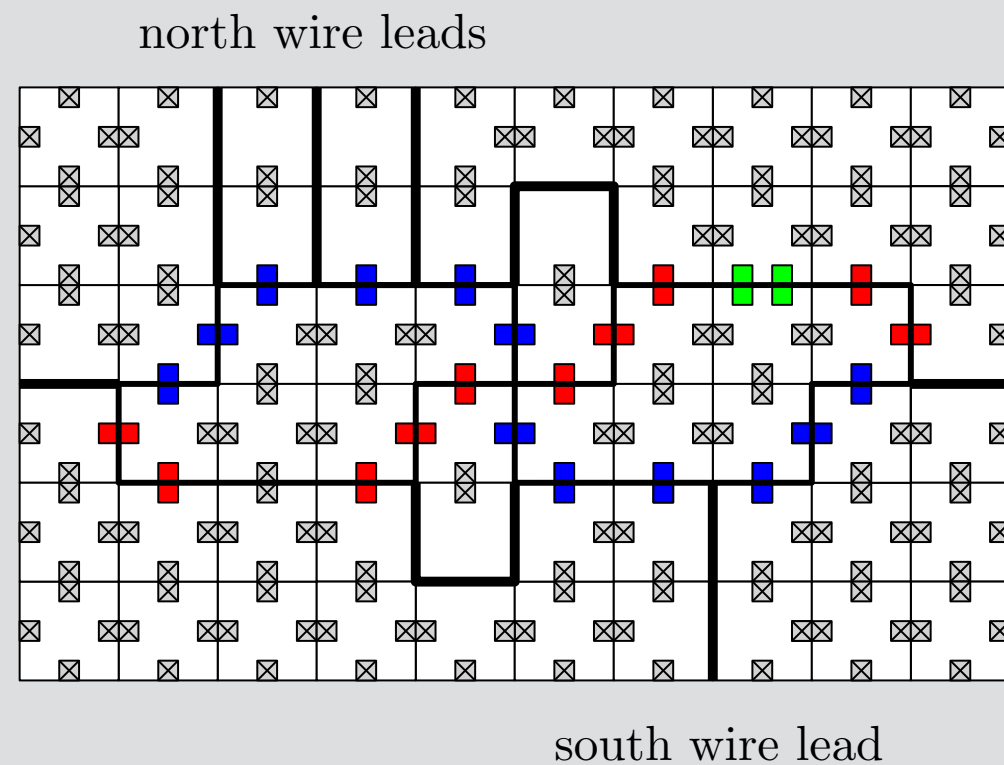
# coNP-hardness Reduction

Reduce from independent set in planar cubic Hamiltonian graphs ([Fleischer et al. 2010])



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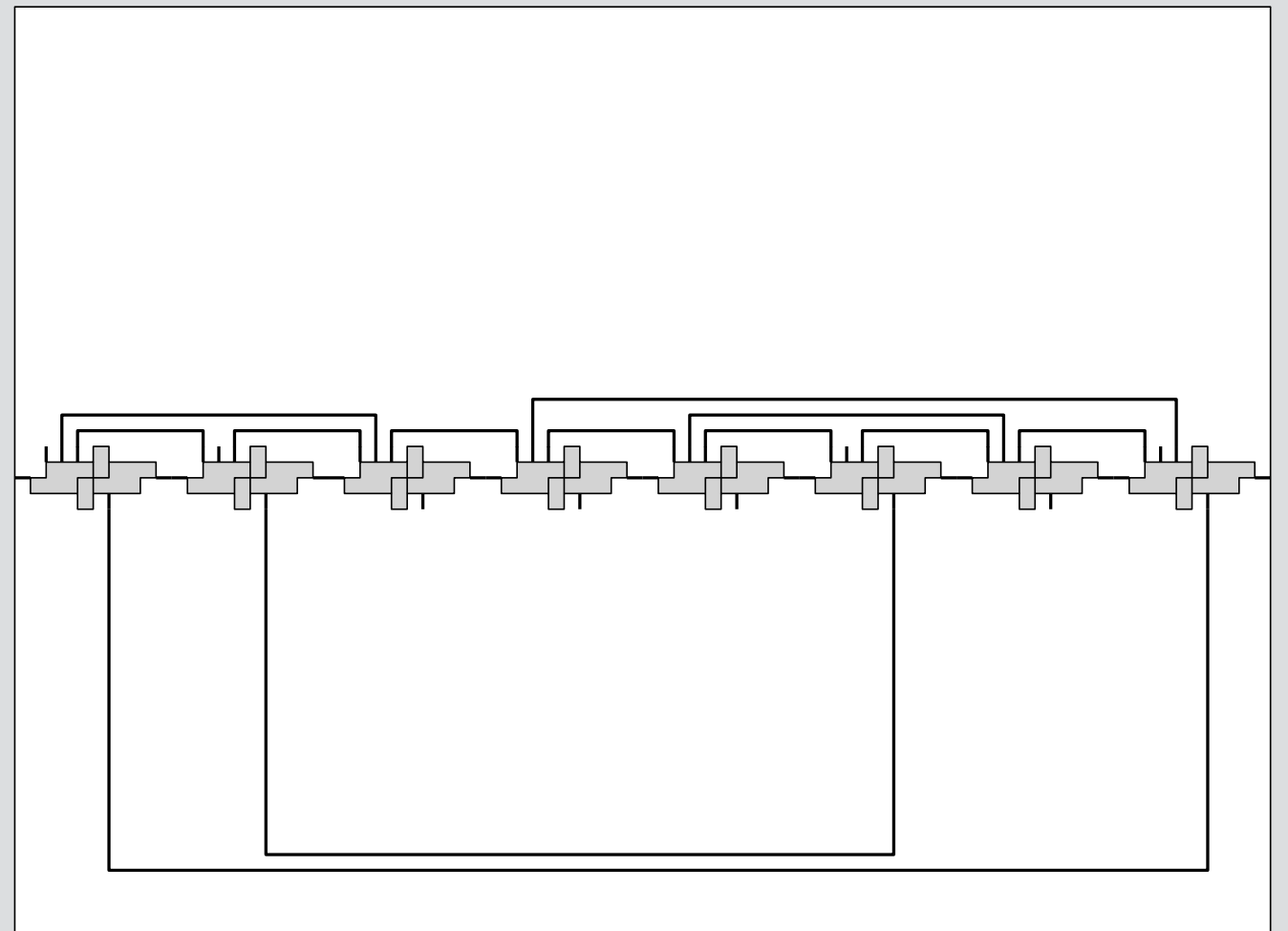
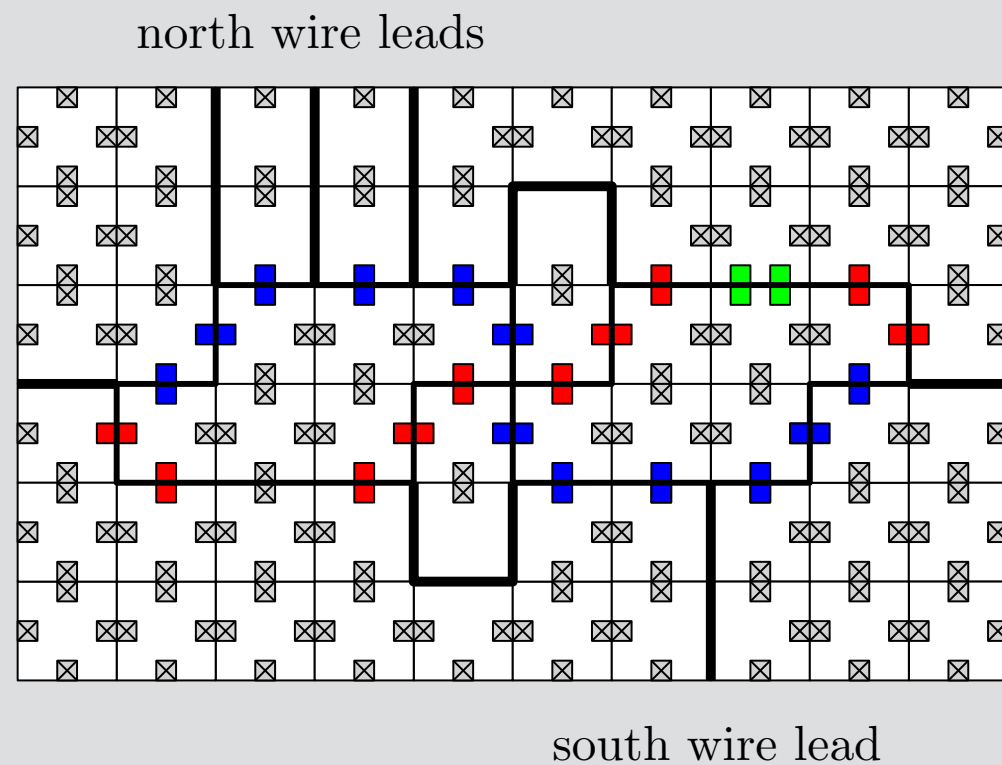
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# coNP-hardness Reduction

Reduce from independent set in planar cubic Hamiltonian graphs ([Fleischer et al. 2010])



$$\tau(n) = \begin{cases} 1 & : n < s/2 \\ 11|V| - k + 1 & : \text{otherwise} \end{cases}$$

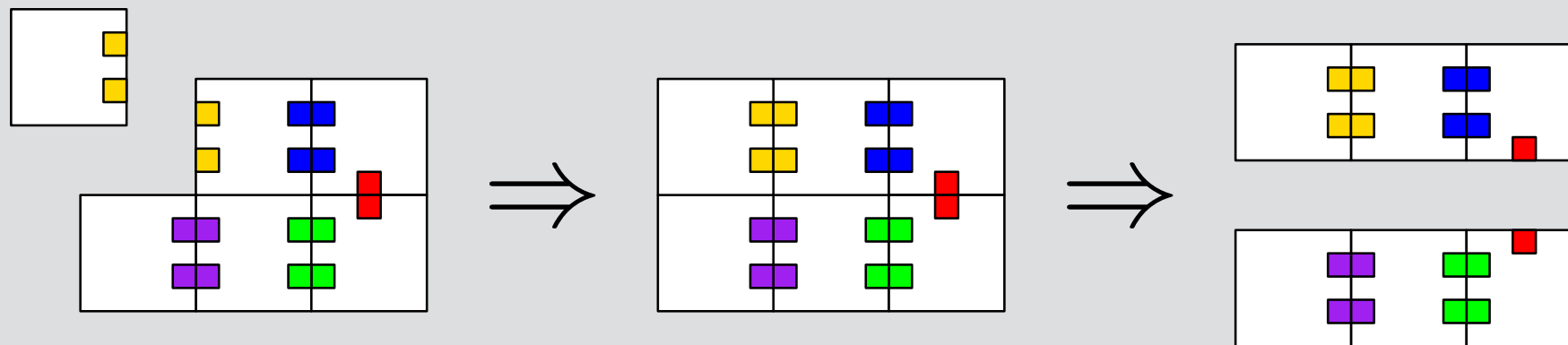
# Conclusion

Temperature functions can yield sophisticated behavior even in simple systems.

Positive and negative: systems are provably more efficient, but (coNP-)harder to design.

Open: positive results with realistic temperature functions. What does “realistic” even mean?

# Size-Dependent Tile Self-Assembly: Constant-Height Rectangles and Stability



Sándor Fekete, Robert Schweller, Andrew Winslow